

## EXPONENTS AND DIAMETERS OF STRONG PRODUCTS OF DIGRAPHS\*

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**Abstract.** The exponent of the strong product of a digraph of order  $m$  and a digraph of order  $n$  is shown to be bounded above by  $m + n - 2$ , with equality for  $Z_m \boxtimes Z_n$ . The exponent and diameter of the strong product of a graph and a digraph are also investigated.

**Key words.** Strong product of digraphs, Exponent, Diameter.

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**1. Introduction.** Let  $D = (V, A)$  be a digraph on  $n$  vertices. Throughout this paper, we assume that  $D$  has no loops and no multiple arcs. A walk from  $u$  to  $v$  in  $D$  is a sequence  $u = u_0, u_1, \dots, u_k = v$  of vertices such that there is an arc from  $u_i$  to  $u_{i+1}$  in  $D$  for each  $i$ . We denote the walk by  $u \rightarrow u_1 \rightarrow u_2 \rightarrow \dots \rightarrow u_{k-1} \rightarrow v$  and its length is  $k$ . We use the notation  $u \xrightarrow{k} v$  when there exists a walk in  $D$  of length  $k$  from  $u$  to  $v$ . The digraph is *primitive* if there is a  $k$  such that  $u \xrightarrow{k} v$  for each pair of vertices  $u$  and  $v$ . Conventionally  $u \xrightarrow{0} u$  is permitted. We say that the smallest such value of  $k$  is the *exponent* of  $D$ , which is denoted by  $\exp(D)$ . Wielandt [9] found that the maximum exponent of a primitive digraph on  $n$  vertices is  $W_n = n^2 - 2n + 2$ . See [1] for more details. Suppose that two digraphs  $D = (V_D, A_D)$  and  $E = (V_E, A_E)$  are given. Let  $V = V_D \times V_E$ . We define

$$A_1 = \{((u_1, u_2), (v_1, v_2)) \in V \times V \mid ((u_1, v_1) \in A_D \text{ and } u_2 = v_2) \\ \text{or } ((u_2, v_2) \in A_E \text{ and } u_1 = v_1)\},$$

and

$$A_2 = \{((u_1, u_2), (v_1, v_2)) \in V \times V \mid (u_1, v_1) \in A_D \text{ and } (u_2, v_2) \in A_E\}.$$

The *strong product*  $D \boxtimes E$  of  $D$  and  $E$  is the digraph  $(V, A_1 \cup A_2)$ . The *Cartesian product*  $D \times E$  and the *direct product*  $D \otimes E$  of  $D$  and  $E$  are defined by  $(V, A_1)$

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and  $(V, A_2)$  respectively. The strong product of graphs is used to define the concept of Shannon capacity which plays an important role in the information theory [8]. Definitions and related results on the product of graphs are provided in [2]. In 1979, Lamprey and Barnes [6] showed that  $\exp(D \times E) \leq (n + m)^2 - 4(n + m) + 5$  for digraphs  $D$  and  $E$  on  $n$  and  $m$  vertices, respectively. They also showed  $\exp(D \otimes E) = \max\{\exp(D), \exp(E)\}$  for primitive digraphs  $D$  and  $E$ . In 1987, Kwasnik [5] studied the exponent of other types of products such as the disjunction and lexicographic products of graphs. Recently, it has been proved in [3] that if  $D$  and  $E$  are digraphs on  $m$  and  $n$  vertices, respectively, and  $D \times E$  is primitive, then  $\exp(D \times E) \leq mn - 1$ . In [3] it was also showed that  $\exp(G \times D) = \exp(G) + \text{diam}(D)$  for a primitive graph  $G$  and a strongly connected bipartite digraph  $D$ , and they computed the exponent of the Cartesian product of two cycles [4]. In this paper, we show

$$\exp(D \boxtimes E) \leq n + m - 2 \quad (1.1)$$

for strongly connected digraphs  $D$  and  $E$  on  $n$  and  $m$  vertices, respectively. Let  $Z_n$  and  $Z_m$  be the directed cycles of order  $n$  and  $m$  respectively. We also prove that

$$\exp(Z_n \boxtimes Z_m) = n + m - 2.$$

As a consequence, the bound in (1.1) is tight. A graph  $G$  is considered as a digraph by treating the edges of  $G$  as bidirectional. In particular, a cycle  $C_n$  of length  $n$  is considered as a digraph in the same manner. For a connected graph  $G$  and a strongly connected digraph  $E$ , we show  $\exp(G \boxtimes E)$  is  $\text{diam}(G \boxtimes E)$  or  $\text{diam}(G \boxtimes E) + 1$ , and we find the condition under which the latter case holds. As a consequence, we compute  $\exp(C_n \boxtimes Z_m)$ .

## 2. Upper bound on the exponent of strong products of two digraphs.

LEMMA 1. Let  $D$  and  $E$  be digraphs,  $u, v \in V_D$  and  $z, w \in V_E$ . If  $u \xrightarrow{t} v$  in  $D$  and  $z \xrightarrow{s} w$  in  $E$ , then  $(u, z) \xrightarrow{\alpha} (v, w)$  in  $D \boxtimes E$  for all  $\alpha$  with  $\max\{t, s\} \leq \alpha \leq t + s$ .

*Proof.* We may assume that  $t \leq s$ . Let  $u \rightarrow u_0 \rightarrow u_1 \rightarrow \cdots \rightarrow u_t = v$  in  $D$  and  $z = z_0 \rightarrow z_1 \rightarrow \cdots \rightarrow z_s = w$  in  $E$ . If  $i = \alpha - s$  for  $0 \leq i \leq t$ , then  $(u, z) = (u_0, z_0) \rightarrow (u_1, z_1) \rightarrow \cdots \rightarrow (u_{t-i}, z_{t-i}) \rightarrow (u_{t-i}, z_{t-i+1}) \rightarrow \cdots \rightarrow (u_{t-i}, z_s) \rightarrow (u_{t-i+1}, z_s) \rightarrow \cdots \rightarrow (u_t, z_s) = (v, w)$  is a walk of length  $t - i + s - (t - i) + t - (t - i) = s + i = \alpha$  in  $D \boxtimes E$ .  $\square$

LEMMA 2. Let  $D$  and  $E$  be strongly connected digraphs,  $u, v \in V_D$  and  $z, w \in V_E$ . If there are a cycle  $C$  passing through  $v$  of length  $k$  in  $D$ ,  $u \xrightarrow{t} v'$  in  $D$  for some vertex  $v'$  of  $C$ , and  $z \xrightarrow{s} w$  in  $E$  for some  $s \geq k - 1$ , then, for all  $\alpha$  with  $\max\{t + k, s\} \leq \alpha$ ,  $(u, z) \xrightarrow{\alpha} (v, w)$  in  $D \boxtimes E$ .

*Proof.* Since  $v$  and  $v'$  are vertices of  $C$ ,  $v' \xrightarrow{l} v$  in  $D$  for some  $l$  with  $0 \leq l \leq k - 1$ . So  $t + l \leq t + k$ . Since  $\alpha - t - l \geq \alpha - t - k \geq 0$ , there is a  $q \geq 0$  such that

$kq \leq \alpha - t - l \leq kq + k - 1$ . Since  $u \xrightarrow{t} v' \xrightarrow{l} v \xrightarrow{kq} v$ ,  $u \xrightarrow{kq+t+l} v$  in  $D$ . Since  $z \xrightarrow{s} w$  in  $E$  and  $\max\{kq + t + l, s\} \leq \alpha \leq kq + t + l + k - 1 \leq kq + t + l + s$ , by Lemma 1,  $(u, z) \xrightarrow{\alpha} (v, w)$  in  $D \boxtimes E$ .  $\square$

**THEOREM 1.** *Let  $D$  and  $E$  be strongly connected digraphs on  $n$  and  $m$  vertices ( $n, m \geq 2$ ), respectively. Then  $D \boxtimes E$  is primitive and*

$$\exp(D \boxtimes E) \leq n + m - 2.$$

*Proof.* It suffices to show that for each pair of vertices  $(u, z), (v, w) \in D \boxtimes E$  and for each  $\alpha \geq n + m - 2$ , we have  $(u, z) \xrightarrow{\alpha} (v, w)$  in  $D \boxtimes E$ . Let  $k$  be the minimum length of the cycles in  $D$  passing through  $v$ , and  $C$  be one such cycle. Let  $t$  be the distance from  $u$  to  $C$  in  $D$ . Then  $k + t \leq n$ .

Let  $l$  be the minimum length of the cycle in  $E$  passing through  $w$  and  $s$  be the distance from  $z$  to  $w$ . Then  $l, s \leq m$ . If  $s \geq k - 1$ , then, by  $\alpha \geq n \geq t + k$  and  $\alpha \geq m \geq s$ , Lemma 2 implies that  $(u, z) \xrightarrow{\alpha} (v, w)$ . If  $s < k - 1$ , then there is a  $q \geq 0$  such that  $lq < k - s - 1 \leq l(q + 1)$ . Then  $l(q + 1) + s = l + lq + s \leq l + k - 2 \leq m + k - 2 \leq n + m - 2 \leq \alpha$ . Since  $z \xrightarrow{s} w \xrightarrow{l(q+1)} w$ ,  $z \xrightarrow{l(q+1)+s} w$ . By Lemma 2,  $(u, z) \xrightarrow{\alpha} (v, w)$ .  $\square$

**THEOREM 2.** *For  $n, m \geq 2$ ,*

$$\exp(Z_n \boxtimes Z_m) = n + m - 2.$$

*Proof.* Let  $Z_n$  be a directed cycle  $v_0 \rightarrow v_1 \rightarrow \cdots \rightarrow v_{n-1} \rightarrow v_0$  and  $Z_m$  be a directed cycle  $w_0 \rightarrow w_1 \rightarrow \cdots \rightarrow w_{m-1} \rightarrow w_0$ . We may assume  $n \leq m$ . Let  $r$  be the residue of  $m - 2$  modulo  $n$ . Suppose  $(v_0, w_0) \xrightarrow{n+m-3} (v_r, w_{n-2})$ . Let  $(v_0, w_0) = (x_0, y_0) \rightarrow (x_1, y_1) \rightarrow \cdots \rightarrow (x_{n+m-3}, y_{n+m-3}) = (v_r, w_{n-2})$  be a path in  $Z_n \boxtimes Z_m$  from  $(v_0, w_0)$  to  $(v_r, w_{n-2})$ . Then there are  $i_0 < i_1 < \cdots < i_s$  and  $j_0 < j_1 < \cdots < j_t$  such that  $i_0 = j_0 = 0$ , for all  $p, q$  with  $0 \leq p \leq s - 1$  and  $0 \leq q \leq t - 1$ ,  $x_{i_p} = x_{i_{p+1}} = \cdots = x_{i_{p+1}-1} \neq x_{i_{p+1}}$ ,  $y_{j_q} = y_{j_{q+1}} = \cdots = y_{j_{q+1}-1} \neq y_{j_{q+1}}$  and  $x_{i_s} = x_{i_{s+1}} = \cdots = x_{n+m-3}$ ,  $y_{j_t} = y_{j_{t+1}} = \cdots = y_{n+m-3}$ . Then for all  $i = 0, 1, \dots, n + m - 3$ ,  $(x_i, y_i) = (x_{i_p}, y_{j_q})$  for some  $p$  and  $q$ . If  $0 \leq i \leq n + m - 4$  and  $(x_i, y_i) = (x_{i_p}, y_{j_q})$ , since  $(x_{i+1}, y_{i+1}) \neq (x_i, y_i)$ ,  $x_{i+1} = x_{i_{p+1}}$  or  $y_{i+1} = y_{j_{q+1}}$ . So if  $(x_i, y_i) = (x_{i_p}, y_{j_q})$ , we can show  $i \leq p + q$ , by induction. If  $x_{i_p} = v_l$  and  $l \neq n - 1$ , since  $(x_{i_p}, x_{i_{p+1}}) \in A_{Z_n}$ ,  $x_{i_{p+1}} = v_{l+1}$ . If  $x_{i_p} = v_{n-1}$ , since  $x_{i_p} = v_{n-1} \rightarrow x_{i_{p+1}}$ ,  $x_{i_{p+1}} = v_0$ . Since  $x_{i_0} = v_0$ , we can show by induction that if  $x_{i_p} = v_l$ ,  $p \equiv l \pmod{n}$ . Similarly, we can show that if  $y_{j_q} = w_k$ ,  $q \equiv k \pmod{m}$ . Since  $y_{j_t} = w_{n-2}$ ,  $t \equiv n - 2 \pmod{m}$ . Since  $t \leq n + m - 3$  and  $n \leq m$ ,  $t = n - 2$ . Since  $x_{i_s} = v_r$ ,  $s \equiv r \pmod{n}$ . So  $s \equiv r \equiv m - 2 \pmod{n}$ .

Since  $s \leq n + m - 3$ ,  $s \leq m - 2$ . So  $n + m - 3 \leq s + t \leq (m - 2) + (n - 2) = n + m - 4$ . This is a contradiction. So  $(v_0, w_0) \xrightarrow{n+m-3} (v_r, w_{n-2})$ . Thus, using Theorem 1,  $\exp(Z_n \boxtimes Z_m) = n + m - 2$ .  $\square$

**3. Exponents and diameters of strong products of digraphs.** For any  $u, v \in V_D$ , the *distance*  $\text{dist}(u, v)$  from  $u$  to  $v$  is the smallest  $k$  such that there is a walk from  $u$  to  $v$  of length  $k$ . The *diameter*  $\text{diam}(D)$  of the strongly connected digraph  $D$  is the maximum of  $\text{dist}(u, v)$  for all  $u, v \in V_D$ .

PROPOSITION 1. *If  $D$  and  $E$  are strongly connected digraphs, then*

$$\text{diam}(D \boxtimes E) = \max\{\text{diam}(D), \text{diam}(E)\}.$$

*Proof.* If  $u, v \in V_D$  and  $z, w \in V_E$ , then by Lemma 1, we have

$$\text{dist}((u, z), (v, w)) \leq \max\{\text{dist}(u, v), \text{dist}(z, w)\}.$$

Thus,  $\text{diam}(D \boxtimes E) \leq \max\{\text{diam}(D), \text{diam}(E)\}$ .

Conversely, if  $u, v \in V_D$ ,  $z, w \in V_E$ , and  $\text{dist}((u, z), (v, w)) = \alpha$ , then  $(u, z) = (u_0, z_0) \rightarrow (u_1, z_1) \rightarrow \cdots \rightarrow (u_\alpha, z_\alpha) = (v, w)$  for some  $(u_i, z_i) \in V_{D \boxtimes E}$  where  $i = 1, 2, \dots, \alpha$ . Thus, there are  $0 = i_0 < i_1 < i_2 < \cdots < i_s \leq \alpha$  such that  $u_{i_p} = u_{i_p+1} = \cdots = u_{i_{p+1}-1} \neq u_{i_{p+1}}$  for all  $p = 0, 1, \dots, s-1$ . Since  $u = u_{i_0} \rightarrow u_{i_1} \rightarrow \cdots \rightarrow u_{i_s} = u_\alpha = v$ ,  $\text{dist}(u, v) \leq s \leq \alpha = \text{dist}((u, z), (v, w))$ . So  $\text{diam}(D) \leq \text{diam}(D \boxtimes E)$ . Similarly,  $\text{diam}(E) \leq \text{diam}(D \boxtimes E)$ . Thus,  $\text{diam}(D \boxtimes E) = \max\{\text{diam}(D), \text{diam}(E)\}$ .  $\square$

LEMMA 3. *Let  $G$  be a connected graph and  $D$  be a strongly connected digraph. If  $(u, z), (v, w) \in V_{G \boxtimes D}$ ,  $((u, z), (v, w)) \in A_{G \boxtimes D}$  and  $z \neq w$ , then  $(u, z) \xrightarrow{k} (v, w)$  in  $G \boxtimes D$  for all  $k \geq 1$ .*

*Proof.* Since  $z \neq w$ ,  $(z, w) \in A_D$ . Since  $G$  is connected, there is  $x \in V_G$  such that  $\{u, x\} \in E_G$ . Since  $(u, z) \rightarrow (x, z) \rightarrow (u, z)$ ,  $(u, z) \xrightarrow{2t} (u, z)$  for all  $t \geq 0$ . If  $u = v$ , since  $(u, z) \xrightarrow{2t} (u, z) \xrightarrow{1} (u, w)$  and  $(u, z) \xrightarrow{2t} (u, z) \xrightarrow{1} (x, z) \xrightarrow{1} (u, w)$ ,  $(u, z) \xrightarrow{2t+1} (u, w) = (v, w)$  and  $(u, z) \xrightarrow{2t+2} (u, w) = (v, w)$  for all  $t \geq 0$ . If  $u \neq v$ ,  $u \rightarrow v$ . Since  $(u, z) \xrightarrow{2t} (u, z) \xrightarrow{1} (v, w)$  and  $(u, z) \xrightarrow{2t} (u, z) \xrightarrow{1} (u, w) \xrightarrow{1} (v, w)$ ,  $(u, z) \xrightarrow{2t+1} (v, w)$  and  $(u, z) \xrightarrow{2t+2} (v, w)$  for all  $t \geq 0$ .  $\square$

THEOREM 3. *If  $G$  is a connected graph and  $D$  is a strongly connected digraph such that  $|V_G| \geq 2$  and  $|V_D| \geq 2$ , then  $\exp(G \boxtimes D)$  is  $\text{diam}(G \boxtimes D)$  or  $\text{diam}(G \boxtimes D) + 1$ . Moreover,  $\exp(G \boxtimes D) = \text{diam}(G \boxtimes D) + 1$  if and only if  $G$  and  $D$  satisfy the following:*

1.  $\text{diam}(D) \geq \text{diam}(G)$ ,

2. there is  $v \in V_D$  such that  $v \xrightarrow{l} v$  for all  $l = 1, 2, \dots, \text{diam}(D)$ ,
3. either  $G$  is not primitive or  $G$  is primitive and  $\exp(G) > \text{diam}(D)$ .

*Proof.* Let  $\text{diam}(G) = m$  and  $\text{diam}(D) = n$ . For all  $u, v \in V_D$ , there is  $u' \in V_D$  such that  $(u, u') \in A_D$ . Since  $\text{diam}(D) = n$ ,  $u' \xrightarrow{t} v$  for some  $t \leq n$ . For all  $x, y \in V_G$ ,  $x \xrightarrow{s} y$  for some  $s \leq m$ . If  $\alpha = \max\{m, n\}$ , by Lemma 1,  $(x, u') \xrightarrow{l} (y, v)$  for some  $l \leq \alpha$ . If  $k \geq \alpha + 1$ , by Lemma 3,  $(x, u) \xrightarrow{k-l} (x, u')$ . Since  $(x, u) \xrightarrow{k-l} (x, u') \xrightarrow{l} (y, v)$ ,  $(x, u) \xrightarrow{k} (y, v)$ . So  $\exp(G \boxtimes D) \leq \alpha + 1 = \text{diam}(G \boxtimes D) + 1$ . Since  $\text{diam}(G \boxtimes D) \leq \exp(G \boxtimes D)$ ,  $\exp(G \boxtimes D)$  is  $\text{diam}(G \boxtimes D)$  or  $\text{diam}(G \boxtimes D) + 1$ .

If  $\exp(G \boxtimes D) = \alpha + 1$ , then there are  $(x, u), (y, v) \in V_{G \boxtimes D}$  such that  $(x, u) \xrightarrow{\alpha} (y, v)$  in  $G \boxtimes D$ . If  $u \xrightarrow{l} v$  for some  $l$  with  $1 \leq l \leq \alpha$ , then there is  $u' \in V_D$  such that  $(u, u') \in A_D$  and  $u' \xrightarrow{l-1} v$ . If  $x \neq y$ , then there is  $x' \in V_G$  such that  $\{x, x'\} \in E_G$  and  $x' \xrightarrow{s-1} y$  where  $s = \text{dist}(x, y)$ . If  $\max\{s, l\} = p$ , then  $p \leq \alpha$ . By Lemmas 1 and 3,  $(x', u') \xrightarrow{p-1} (y, v)$  and  $(x, u) \xrightarrow{\alpha-p+1} (x', u')$ . So  $(x, u) \xrightarrow{\alpha} (y, v)$ . This is a contradiction. If  $x = y$ , by Lemma 3,  $(x, u) \xrightarrow{\alpha-l+1} (x, u')$ . Since  $(x, u) \xrightarrow{\alpha-l+1} (x, u') \xrightarrow{l-1} (x, v)$ ,  $(x, u) \xrightarrow{\alpha} (x, v) = (y, v)$ . This is a contradiction. So  $u \xrightarrow{l} v$  for all  $l$  such that  $1 \leq l \leq \alpha$ . If  $u \neq v$ , let  $d = \text{dist}(u, v)$ . Then  $1 \leq d \leq \alpha$  and  $u \xrightarrow{d} v$ . This is a contradiction. So  $u = v$  and  $G$  and  $D$  satisfy condition (2).

Since  $D$  is strongly connected, there is  $\tilde{u} \in V_D$  such that  $(\tilde{u}, u) \in A_D$ . If  $\text{dist}(u, \tilde{u}) = r \leq \alpha$ , since  $u \xrightarrow{r+1} u$ ,  $r + 1 \geq \alpha + 1$ . Since  $r \leq n \leq \alpha$ ,  $r = n = \alpha$ . So  $n \geq m$ . Thus,  $G$  and  $D$  satisfy condition (1). Since  $(x, u) \xrightarrow{\alpha} (y, u)$ ,  $x \xrightarrow{\alpha} y$  in  $G$ . So  $G$  and  $D$  satisfy condition (3).

Conversely, if  $G$  and  $D$  satisfy conditions (1), (2) and (3), then there is  $u \in V_D$  such that  $u \xrightarrow{l} u$  for all  $l$  such that  $1 \leq l \leq n = \alpha$ , and there are  $x, y \in V_G$  such that  $x \xrightarrow{\alpha} y$  in  $G$ . If  $(x, u) \xrightarrow{\alpha} (y, u)$ , then  $(x, u) = (x_0, u_0) \rightarrow (x_1, u_1) \rightarrow \dots \rightarrow (x_\alpha, u_\alpha) = (y, u)$  for some  $x_0, x_1, \dots, x_\alpha \in V_G$  and  $u_0, u_1, \dots, u_\alpha \in V_D$ . If  $u_0 = u_1 = \dots = u_\alpha = u$ , since  $x = x_0 \rightarrow x_1 \rightarrow \dots \rightarrow x_\alpha = y$ ,  $x \xrightarrow{\alpha} y$ . This is a contradiction. If  $u_i \neq u$  for some  $i$ , there are  $0 = i_0 < i_1 < \dots < i_s \leq \alpha$  such that  $u_{i_p} = u_{i_p+1} = \dots = u_{i_{p+1}-1} \neq u_{i_{p+1}}$  for all  $p = 0, 1, \dots, s-1$  and  $u_{i_s} = u_{i_s+1} = \dots = u_\alpha$ . Since  $u = u_{i_0} \rightarrow u_{i_1} \rightarrow \dots \rightarrow u_{i_s} \rightarrow u_\alpha = u$ ,  $u \xrightarrow{s} u$ . Since  $u_i \neq u$  for some  $i$ ,  $1 \leq s \leq \alpha$ . This is a contradiction. So  $(x, u) \xrightarrow{\alpha} (y, u)$ . Thus,  $\exp(G \boxtimes D) = \alpha + 1$ .  $\square$

COROLLARY 1. If  $G$  and  $H$  are connected graphs, then

$$\exp(G \boxtimes H) = \text{diam}(G \boxtimes H)$$

except when both  $G$  and  $H$  are complete graphs.

*Proof.* If  $\exp(G \boxtimes H) = \text{diam}(G \boxtimes H) + 1$ , since  $v \xrightarrow{2} v$  for all  $v \in V_H$ , by Theorem 3,  $1 \leq \text{diam}(G) \leq \text{diam}(H) = 1$ . So  $G$  and  $H$  are complete graphs.  $\square$

Note that the strong product of two complete graphs is also a complete graph, whose exponent is 2.

COROLLARY 2.

$$\exp(C_n \boxtimes Z_m) = \begin{cases} \lfloor \frac{n}{2} \rfloor, & \text{if } n \geq 2m \\ m-1, & \text{if } n \text{ is odd and } n \leq m \\ m, & \text{if } n \text{ is even and } n \leq 2m-2, \text{ or } n \text{ is odd} \\ & \text{and } m+1 \leq n \leq 2m-1. \end{cases}$$

*Proof.* If  $\exp(C_n \boxtimes Z_m) = \text{diam}(C_n \boxtimes Z_m) + 1$ ,  $\text{diam}(C_n) = \lfloor \frac{n}{2} \rfloor \leq \text{diam}(Z_m) = m-1$ . So  $n \leq 2m-1$ . Moreover,  $C_n$  is not primitive, or  $C_n$  is primitive and  $\exp(C_n) = n-1 > \text{diam}(Z_m) = m-1$ . So  $n$  is even, or  $n$  is odd and  $n \geq m+1$ . Thus, if  $n$  is even,  $n \leq 2m-2$ . And if  $n$  is odd,  $m+1 \leq n \leq 2m-1$ . In this case,  $\exp(C_n \boxtimes Z_m) = (m-1) + 1 = m$ . Otherwise,  $\exp(C_n \boxtimes Z_m) = \text{diam}(C_n \boxtimes Z_m) = \max\{\lfloor \frac{n}{2} \rfloor, m-1\} = \begin{cases} \lfloor \frac{n}{2} \rfloor, & n \geq 2m \\ m-1, & \text{if } n \leq 2m-1. \end{cases} \quad \square$

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