

THE MINIMUM RANK PROBLEM OVER FINITE FIELDS*

JASON GROUT^{\dagger}

Abstract. The structure of all graphs having minimum rank at most k over a finite field with q elements is characterized for any possible k and q. A strong connection between this characterization and polarities of projective geometries is explained. Using this connection, a few results in the minimum rank problem are derived by applying some known results from projective geometry.

Key words. Minimum rank, Symmetric matrix, Finite field, Projective geometry, Polarity graph, Bilinear symmetric form

AMS subject classifications. 05C50, 05C75, 15A03, 05B25, 51E20.

1. Introduction. Given a field F and a simple undirected graph G on n vertices (i.e., an undirected graph without loops or multiple edges), let S(F,G) be the set of symmetric $n \times n$ matrices A with entries in F satisfying $a_{ij} \neq 0, i \neq j$, if and only if ij is an edge in G. There is no restriction on the diagonal entries of the matrices in S(F,G). Let

 $mr(F,G) = \min\{\operatorname{rank} A \mid A \in S(F,G)\}.$

Let $\mathcal{G}_k(F) = \{G \mid \operatorname{mr}(F,G) \leq k\}$, the set of simple graphs with minimum rank at most k.

The problem of finding mr(F, G) and describing $\mathcal{G}_k(F)$ has recently attracted considerable attention, particularly for the case in which $F = \mathbb{R}$ (see [29, 17, 26, 25, 27, 13, 33, 5, 9, 22, 2, 11, 6, 7, 10, 18, 4]). The minimum rank problem over \mathbb{R} is a sub-problem of a much more general problem, the inverse eigenvalue problem for symmetric matrices: given a family of real numbers, find every symmetric matrix that has the family as its eigenvalues. More particularly, the minimum rank problem is a sub-problem of the inverse eigenvalue problem for graphs, which fixes a zero/nonzero pattern for the symmetric matrices considered in the inverse eigenvalue problem. The minimum rank problem can also be thought of in this way: given a fixed pattern of off-diagonal zeros, what is the smallest rank that a symmetric matrix having that pattern can achieve?

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[†] Department of Mathematics and Computer Science, Drake University, 2507 University Avenue, Des Moines, IA 50311, USA (jason.grout@drake.edu).



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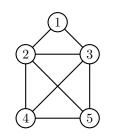


Fig. 1.1: A labeled full house graph

Up to the addition of isolated vertices, it is easy to see that $\mathcal{G}_1(F) = \{K_n \mid n \in \mathbb{N}\}$ for any field F. In [9] and [10], $\mathcal{G}_2(F)$ was characterized for any field F both in terms of forbidden subgraphs and in terms of the structure of the graph complements. The forbidden subgraph characterizations in these papers used ten or fewer graphs for each value of k. Restricting our focus to finite fields, let \mathbb{F}_q denote the finite field with q elements. Ding and Kotlov [18] independently used structures similar to the structures that we use in this paper to obtain some special cases of some structural results in this paper, as well as an upper bound for the sizes of minimal forbidden subgraphs characterizing $\mathcal{G}_k(\mathbb{F}_q)$ for any k and any q. The latter result implies that there are a finite number of forbidden subgraphs characterizing $\mathcal{G}_k(\mathbb{F}_q)$. For example, in [8], $\mathcal{G}_3(\mathbb{F}_2)$ was characterized by 62 forbidden subgraphs. This characterization and further computations confirm our intuition that the forbidden subgraph characterizations of $\mathcal{G}_k(\mathbb{F}_q)$ quickly become complicated as k increases.

In this paper, we will characterize the structure of graphs in $\mathcal{G}_k(\mathbb{F}_q)$ for any k and any q. The characterization is simply stated and has a very strong connection to projective geometry over finite fields. At the end of the paper, we will list a few of the ramifications of this connection to projective geometry.

We adopt the following notation dealing with fields, vector spaces, and matrices. Given a field F, the group of nonzero elements under multiplication is denoted F^{\times} and the vector space of dimension k over F is denoted F^k . Given a matrix M, the principal submatrix lying in the rows and columns x_1, x_2, \ldots, x_m is denoted $M[x_1, x_2, \ldots, x_m]$.

To illustrate the field-dependence of minimum rank, we recall from [10] the full house graph in Figure 1.1 (there called $(P_3 \cup 2K_1)^c$), which is the only graph on 5 or fewer vertices for which the minimum rank is field-dependent.



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If $F \neq \mathbb{F}_2$, there are elements $a, b \neq 0$ in F such that $a + b \neq 0$. Then

$$\begin{bmatrix} a & a & a & 0 & 0 \\ a & a+b & a+b & b & b \\ a & a+b & a+b & b & b \\ 0 & b & b & b & b \\ 0 & b & b & b & b \end{bmatrix} \in S(F, \text{full house}),$$

which shows that mr(F, full house) = 2. The case $F = \mathbb{F}_2$ gives a different result. Let A be any matrix in $S(\mathbb{F}_2, full house)$. Then for some $d_1, d_2, \ldots, d_5 \in \mathbb{F}_2$,

$$A = \begin{vmatrix} d_1 & 1 & 1 & 0 & 0 \\ 1 & d_2 & 1 & 1 & 1 \\ 1 & 1 & d_3 & 1 & 1 \\ 0 & 1 & 1 & d_4 & 1 \\ 0 & 1 & 1 & 1 & d_5 \end{vmatrix} \quad \text{and} \quad \det(A[\{1,2,5\},\{1,3,4\}]) = \begin{vmatrix} d_1 & 1 & 0 \\ 1 & 1 & 1 \\ 0 & 1 & 1 \end{vmatrix} = 1,$$

where $A[\{1, 2, 5\}, \{1, 3, 4\}]$ is the submatrix of A lying in rows $\{1, 2, 5\}$ and columns $\{1, 3, 4\}$. Therefore $mr(\mathbb{F}_2, \text{full house}) \geq 3$. Setting each d_i to 1 verifies the statement that $mr(\mathbb{F}_2, \text{full house}) = 3$.

In spite of this dependence on the field, there are a number of results about minimum rank that are field independent. For example, the minimum rank of a tree is field independent (see any of [3], [31], or [14]). Many of the forbidden subgraphs classifying $\mathcal{G}_3(\mathbb{F}_2)$ that are found in [8] are also forbidden subgraphs for $\mathcal{G}_3(F)$ for any field F. These results and others demonstrate that results obtained over finite fields can provide important insights for other fields.

The presentation of material in this paper is oriented towards a reader that is familiar with concepts from linear algebra and graph theory. In the rest of this section, we will review some of our conventions in terminology from graph theory.

In this paper, graphs are undirected, may have loops, but will not have multiple edges between vertices. To simplify our drawings, a vertex with a loop (a *looped vertex*) will be filled (black) and a vertex without a loop (a *nonlooped vertex*) will be empty (white). A simple graph is a graph without loops. Let G be a graph with some loops and \hat{G} be the simple version of G obtained by deleting all loops. We say that a matrix in $S(F, \hat{G})$ corresponds to the simple graph \hat{G} . A matrix $A \in S(F, \hat{G})$ corresponds to G if a_{ii} is nonzero exactly when the vertex *i* has a loop in G. Note that if a matrix corresponds to a looped graph, then it also corresponds to the simple version of the graph.

We recall some notation from graph theory.

DEFINITION 1.1. Given two graphs G and H with disjoint vertex sets V(G) and V(H) and edge sets E(G) and E(H), the union of G and H, denoted $G \cup H$, has

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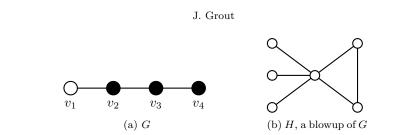


Fig. 1.2: Graphs in Example 1.5

vertices $V(G) \cup V(H)$ and edges $E(G) \cup E(H)$. The *join* of G and H, denoted $G \vee H$, has vertices $V(G) \cup V(H)$ and edges $E(G) \cup E(H) \cup \{uv \mid u \in V(G), v \in V(H)\}$. The complement of the graph G, denoted G^c , has vertices V(G) and edges $\{uv \mid u, v \in$ $V(G), uv \notin E(G)\}$. Note that a vertex is looped in G if and only if it is nonlooped in G^c .

DEFINITION 1.2. The simple complete graph on n vertices will be denoted by K_n and has vertices $\{1, 2, ..., n\}$ and edges $\{xy \mid x, y \in V(K_n), x \neq y\}$. The simple complete multipartite graph $K_{s_1,s_2,...,s_m}$ is defined as $K_{s_1}^c \vee K_{s_2}^c \vee \cdots \vee K_{s_m}^c$.

DEFINITION 1.3. Two vertices in a graph are *adjacent* if an edge connects them. A *clique* in a graph is a set of pairwise adjacent vertices. An *independent set* in a graph is a set of pairwise nonadjacent vertices.

The next definition extends a standard definition introduced in [28] and is used in random graph theory in connection with the regularity lemma.

DEFINITION 1.4. A blowup of a graph G with vertices $\{v_1, v_2, \ldots, v_n\}$ is a new simple graph H constructed by replacing each nonlooped vertex v_i in G with a (possibly empty) independent set V_i , each looped vertex v_i with a (possibly empty) clique V_i , and each edge $v_i v_j$ in G $(i \neq j)$ with the edges $\{xy \mid x \in V_i, y \in V_j\}$ in H.

EXAMPLE 1.5. Let G be the graph labeled in Figure 1.2(a).

Let $|V_1| = 3$, $|V_2| = 1$, $|V_3| = 2$, and $|V_4| = 0$. Then we obtain the simple blowup graph H in Figure 1.2(b). It is useful to see how matrices corresponding to a graph and a blowup of the graph are related. Over \mathbb{F}_3 , let

$$M = \begin{bmatrix} 0 & 2 & 0 & 0 \\ \hline 2 & 1 & 1 & 0 \\ \hline 0 & 1 & 1 & 1 \\ \hline 0 & 0 & 1 & 1 \end{bmatrix} \text{ and } N = \begin{bmatrix} 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 2 & 0 & 0 \\ \hline 0 & 0 & 0 & 1 & 0 & 0 \\ \hline 1 & 2 & 1 & 0 & 1 & 1 \\ \hline 0 & 0 & 0 & 1 & 0 & 1 \\ \hline 0 & 0 & 0 & 1 & 1 & 2 \end{bmatrix}$$



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Fig. 2.1: Graphs in Theorem 2.2

Then M is an example of a matrix corresponding to G and N is an example of a matrix corresponding to H. Note that, for example, the entry m_{11} was replaced with a 3×3 zero block in N, the entry m_{12} was replaced with a 3×1 nonzero block in N, the entry m_{12} was replaced with a 3×1 nonzero block in N, the entries in the last row and column of M were replaced with empty blocks (i.e., erased), and the diagonal entries of N were changed to whatever was desired. These substitutions of block matrices correspond to the vertex substitutions used to construct H.

We will introduce our method by presenting a proof of a special case of a characterization theorem from [10] which characterizes $\mathcal{G}_2(\mathbb{F}_2)$. We will then generalize this proof into a characterization of all simple graphs in $\mathcal{G}_k(\mathbb{F}_q)$ for any k and q. After giving examples for some specific k and q, we will describe the strong connection to projective geometry and list some consequences of this connection.

2. A new approach to a recent result. We will introduce our method by giving a proof of a special case of Theorems 5 and 6 of [10].

THEOREM 2.1 ([10]). Let G be a simple graph on n vertices. Then $mr(\mathbb{F}_2, G) \leq 2$ if and only if the simple version of G^c is either of the form

$$(K_{s_1} \cup K_{p_1,q_1}) \vee K_r$$

for some appropriate nonnegative integers s_1 , p_1 , q_1 , and r, or of the form

$$(K_{s_1} \cup K_{s_2} \cup K_{s_3}) \vee K_r$$

for some appropriate nonnegative integers s_1 , s_2 , s_3 , and r.

We first rephrase Theorem 2.1 using blowup graph terminology.

THEOREM 2.2 ([10]). Let G be a simple graph on n vertices. Then $mr(\mathbb{F}_2, G) \leq 2$ (i.e., $G \in \mathcal{G}_2(\mathbb{F}_2)$) if and only if G is a blowup of either of the graphs in Figure 2.1.

In the proof of this result, we will need the following lemma and corollary, which hold in any field. We will then give a proof of Theorem 2.2.



LEMMA 2.3 ([15, Theorem 8.9.1]). Let A be an $n \times n$ symmetric matrix of rank k. Then there is an invertible principal $k \times k$ submatrix B of A and a $k \times n$ matrix U such that

$$A = U^t B U_t$$

COROLLARY 2.4. Let A be an $n \times n$ symmetric matrix. Then rank $A \leq k$ if and only if there is some invertible $k \times k$ matrix B and $k \times n$ matrix U such that $A = U^t B U$.

Proof. Let A have rank $r \leq k$. Then by Lemma 2.3, there is an invertible $r \times r$ matrix B_1 and an $r \times n$ matrix U_1 such that $A = U_1^t B_1 U_1$. Let $B_2 = \begin{bmatrix} B_1 & O \\ O & I_{k-r} \end{bmatrix}$ and $U_2 = \begin{bmatrix} U_1 \\ O \end{bmatrix}$ (where O represents a zero matrix of the appropriate size). Then $A = U_2^t B_2 U_2$. The reverse implication follows from the rank inequality rank $(U^t B U) \leq$ rank B. \Box

Recall that two square matrices A and B are congruent if there exists some invertible matrix C such that $A = C^t B C$. It is straightforward to show that congruence is an equivalence relation. Let \mathcal{B} consist of one representative from each congruence equivalence class of invertible symmetric $k \times k$ matrices. By Corollary 2.4, if A is a symmetric $n \times n$ matrix with rank $A \leq k$, then $A \in \{U^t B U \mid B \in \mathcal{B}, U \mid k \times n \text{ matrix}\}$.

We now proceed with the proof of Theorem 2.2.

Proof. [Proof of Theorem 2.2] First, we compute a suitable \mathcal{B} , a set of representatives from the congruence classes of invertible symmetric 2×2 matrices over \mathbb{F}_2 . If an invertible symmetric 2×2 matrix B over \mathbb{F}_2 has a nonzero diagonal entry, then $B = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}$, $B = \begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix}$, or $B = I_2$. In any of these three cases, $B^t B B = I_2$, so B is congruent to the identity matrix I_2 . If an invertible symmetric 2×2 matrix B over \mathbb{F}_2 has all zeros on the diagonal, then the off-diagonal entries must be nonzero, so $B = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$. In this case,

$$\begin{bmatrix} a & c \\ b & d \end{bmatrix} \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} a & b \\ c & d \end{bmatrix} = \begin{bmatrix} ac+ac & ad+bc \\ ad+bc & bd+bd \end{bmatrix} = \begin{bmatrix} 0 & ad+bc \\ ad+bc & 0 \end{bmatrix},$$

so any matrix congruent to B will have a zero diagonal. Therefore, a suitable \mathcal{B} is

$$\mathcal{B} = \left\{ I_2, \left[\begin{array}{cc} 0 & 1 \\ 1 & 0 \end{array} \right] \right\}.$$



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Because U is a matrix with entries in \mathbb{F}_2 , the columns of U are members of the finite set

$$\left[\left[\begin{array}{c} 1\\0 \end{array} \right], \left[\begin{array}{c} 0\\1 \end{array} \right], \left[\begin{array}{c} 1\\1 \end{array} \right], \left[\begin{array}{c} 0\\0 \end{array} \right] \right\}.$$

Let A be a symmetric $k \times k$ matrix. For any $n \times n$ permutation matrix P, the graphs of A and P^tAP are isomorphic. Therefore we may assume that identical columns of U are contiguous and write $U = \begin{bmatrix} E_1 & E_2 & J & O \end{bmatrix}$ where E_1 is $2 \times p$ matrix with each column equal to $\begin{bmatrix} 1 \\ 0 \end{bmatrix}$, E_2 is $2 \times q$ matrix with each column equal to $\begin{bmatrix} 0 \\ 1 \end{bmatrix}$, J is a $2 \times r$ matrix with each entry equal to 1, and O is a $2 \times t$ zero matrix. Then either

$$A = \begin{bmatrix} E_1^T \\ E_2^T \\ J^T \\ O^T \end{bmatrix} \begin{bmatrix} E_1 & E_2 & J & O \end{bmatrix} = \begin{bmatrix} J_p & O & J_{p,r} & O \\ O & J_q & J_{q,r} & O \\ J_{r,p} & J_{r,q} & O_r & O \\ O & O & O & O_t \end{bmatrix}$$

or else

$$A = \begin{bmatrix} E_1^{\mathrm{T}} \\ E_2^{\mathrm{T}} \\ J^{\mathrm{T}} \\ O^{\mathrm{T}} \end{bmatrix} \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} E_1 & E_2 & J & O \end{bmatrix} = \begin{bmatrix} O_p & J_{p,q} & J_{p,r} & O \\ J_{q,p} & O_q & J_{q,r} & O \\ J_{r,p} & J_{r,q} & O_r & O \\ O & O & O & O_t \end{bmatrix},$$

where J is an all-ones matrix, O is a zero matrix, and subscripts of J and O denote the dimensions of the matrix.

Any simple graph corresponding to the first matrix is a blowup of the graph in Figure 2.1(a), while any simple graph corresponding to the second matrix is a blowup of the graph in Figure 2.1(b). Thus we have established Theorem 2.2. \Box

OBSERVATION 2.5. Note that every block in the above matrices is either a O matrix or a J matrix. Consequently, we could have obtained the zero/nonzero form of the matrices with rank at most 2 by only considering $U = \begin{bmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 \end{bmatrix}$ and computing

$$A = U^{t}U = \begin{bmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

and

$$A = U^{t}B_{2}U = \begin{bmatrix} 1 & 0 \\ 0 & 1 \\ 1 & 1 \\ 0 & 0 \end{bmatrix} \begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \end{bmatrix} = \begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}.$$



The nonzero diagonal entries correspond to loops in our graphs. This simplified procedure again yields the graphs in Figure 2.1.

In the proof of Theorem 2.2, we noted that any U could be written in a standard form. In Observation 2.5, we saw how the standard form of U could be simplified to take advantage of the theorem being about blowup graphs. We will now discuss the reasoning behind these constructions and show that an analogous standard form of U exists for any finite field and any k.

Because we construct the graphs using representatives of congruence classes, it is important for any simplified U to have the property that if B and \hat{B} are congruent, then $U^t B U$ and $U^t \hat{B} U$ correspond to isomorphic graphs. The following lemma shows that if we take a matrix U where the columns consist of all vectors in \mathbb{F}_q^k , like in Observation 2.5, and if B and \hat{B} are congruent, then $U^t B U$ and $U^t \hat{B} U$ correspond to isomorphic graphs.

LEMMA 2.6. Let U be the matrix with columns $\{v \mid v \in \mathbb{F}_q^k\}$. Let B and C be invertible $k \times k$ matrices with B symmetric. Then the graphs corresponding to U^tBU and $U^t(C^tBC)U$ are isomorphic.

Proof. Since every vector in \mathbb{F}_q^k appears as a column of U and the mapping $x \mapsto Cx$ is one-to-one, CU is just a column permutation of U. This permutation induces a relabeling of the graph $U^t B U$ to give the graph of $(CU)^t B(CU) = U^t (C^t B C) U$. \Box

Though this invariance property with respect to congruent matrices does not hold for an arbitrary U, there is another smaller U which does have the same property. We first need some preliminary material. Then we will introduce this new U in Lemma 2.9.

DEFINITION 2.7. Let F be a field. Two nonzero vectors $v_1, v_2 \in F^k$ are projectively equivalent if there exists some nonzero $c \in F$ such that $v_1 = cv_2$.

It is easy to check that projective equivalence is in fact an equivalence relation on the vectors in V.

We pause to note that replacing a column of U with a projectively equivalent column does not affect the graph corresponding to $U^t B U$. To see this, let $U = [u_1 \ u_2 \ \cdots \ u_n]$ and let $i \in \{1, 2, \ldots, n\}$. Let \hat{U} be the matrix obtained from U by replacing the column u_i with cu_i for some nonzero $c \in F$. Then the i, j entry of $\hat{U}^t B \hat{U}, (cu_i)^t B u_j$ if $i \neq j$ or $(cu_i)^t B (cu_i)$ if i = j, is zero if and only if the i, j entry of $U^t B U, u_i^t B u_j$, is zero. Thus the graphs associated with $U^t B U$ and $\hat{U}^t B \hat{U}$ are equal.

LEMMA 2.8. Let F be any field, let $x \in F^k$, let \bar{x} denote the projective equivalence class of x, and let $P = \bigcup_{x \in F^k - \bar{0}} \{\bar{x}\}$, the set of projective equivalence classes in F^k . Let C be an invertible matrix. Then the map $f: P \to P$ defined by $f: \bar{x} \mapsto \overline{Cx}$ is a



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bijection.

Proof. The function f is well-defined since if Cx = y, then for any nonzero $k \in F$, $\overline{C(kx)} = \overline{kCx} = \overline{ky} = \overline{y}$. If $\overline{Cx_1} = \overline{Cx_2}$, then for some nonzero $k \in F$, $kCx_1 = Cx_2$, which implies $C(kx_1 - x_2) = 0$, giving $kx_1 = x_2$ since C is invertible. Therefore $\overline{x_1} = \overline{x_2}$ and f is injective. Surjectivity of f also follows from the hypothesis that C is invertible. \Box

LEMMA 2.9. Let $\overline{x_1}, \overline{x_2}, \ldots, \overline{x_m}$ be the projective equivalence classes of $\mathbb{F}_q^k - \vec{0}$, with each x_i as a chosen representative from its class. Let $U = [x_1 \ x_2 \ \cdots \ x_m]$, the matrix with column vectors x_1, x_2, \ldots, x_m . Let B and C be invertible $k \times k$ matrices with B symmetric. Then the graphs corresponding to U^tBU and $U^t(C^tBC)U$ are isomorphic.

Proof. Let T = CU. Denote the *i*th column of U by u_i and the *i*th column of T by t_i . By Lemma 2.8, the sequence of projective equivalence classes $\overline{t_1}, \overline{t_2}, \ldots, \overline{t_n}$ is just a permutation of the sequence $\overline{u_1}, \overline{u_2}, \ldots, \overline{u_n}$. Form the matrix S in which the *i*th column, s_i , is u_j if $\overline{t_i} = \overline{u_j}$, so that S is a column permutation of U and $\overline{s_i} = \overline{t_i}$. Then the graph corresponding to $U^t(C^tBC)U = (CU)^tB(CU) = T^tBT$ is isomorphic to the graph corresponding to S^tBS by the reasoning preceding Lemma 2.8, which is in turn just a relabeling of the graph corresponding to U^tBU . \Box

We now find a standard form for any matrix U, as in our proof of Theorem 2.2. Let U be a $k \times n$ matrix over \mathbb{F}_q and let B be an invertible symmetric $k \times k$ matrix over \mathbb{F}_q . Let $\overline{x_1}, \overline{x_2}, \ldots, \overline{x_m}$ be the projective equivalence classes of $\mathbb{F}_q^k - \vec{0}$, with each x_i as a chosen representative from its class. For each nonzero column u_i of U, replace u_i with the chosen representative of $\overline{u_i}$. Then permute the columns of U so that the matrix is of the form $\hat{U} = [X_1 \ X_2 \ \cdots \ X_m \ O]$, where each X_i is a block matrix of columns equal to x_i and O is a zero block matrix. Note that some of these blocks may be empty. Let G be the simple graph corresponding to $U^t BU$ and let \hat{G} be the simple graph corresponding to $\hat{U}^t B\hat{U}$. From our results above, G is isomorphic to \hat{G} .

As illustrated in Observation 2.5, we can obtain the zero/nonzero structure of the block matrix $\hat{U}^t B \hat{U}$ by simply deleting all duplicate columns of \hat{U} . Deleting these duplicate columns of \hat{U} leaves a matrix that can be obtained from $\tilde{U} = [x_1 \ x_2 \ \cdots \ x_m \ 0]$ by deleting the columns of \tilde{U} corresponding to empty blocks of \hat{U} . Let \tilde{G} be the (looped) graph corresponding to $\tilde{U}^t B \tilde{U}$. Then \hat{G} is a blowup of \tilde{G} , which implies that G is a blowup of \tilde{G} .

Furthermore, let \mathcal{B} be a set consisting of one representative from each congruence class of invertible symmetric $k \times k$ matrices and let \hat{B} be the representative that is congruent to B. Then from Lemma 2.9, the graphs corresponding to $\tilde{U}^t B \tilde{U}$ and $\tilde{U}^t B \tilde{U}$ are isomorphic.



There is another simplification we can make. Notice that both graphs displayed in Theorem 2.2 have an isolated nonlooped vertex. This vertex came from the zero column vectors in U and corresponds to the fact that adding any number of isolated vertices to a graph does not change its minimum rank. In any theorem like Theorem 2.2, each graph from which we construct blowups will always have this isolated nonlooped vertex and so will be of the form $G \cup K_1$. Note that in constructing such a graph G, it is enough to assume that \tilde{U} in the above paragraphs does not have a zero column vector.

DEFINITION 2.10. Let $\overline{x_1}, \overline{x_2}, \ldots, \overline{x_m}$ be the projective equivalence classes of $\mathbb{F}_q^k - \vec{0}$, with each x_i as a chosen representative from its class. Let \mathcal{B} be a set consisting of one representative from each congruence class of invertible symmetric $k \times k$ matrices. Let $U = [x_1 \ x_2 \ \cdots \ x_m]$, the matrix with column vectors x_1, x_2, \ldots, x_m . We define the set of graphs $\mathfrak{g}_k(\mathbb{F}_q)$ as the set of graphs corresponding to the matrices in $\{U^t B U \mid B \in \mathcal{B}\}$.

We now have the following result (recall that K_1 has no loop).

THEOREM 2.11. A simple graph G is in $\mathcal{G}_k(\mathbb{F}_q)$ if and only if G is a blowup of some graph in $\{H \cup K_1 \mid H \in \mathfrak{g}_k(\mathbb{F}_q)\}$.

Proof. Let G be a simple graph in $\mathcal{G}_k(\mathbb{F}_q)$. Let $A \in S(\mathbb{F}_q, G)$ be a matrix with rank $A \leq k$. Then $A = U^t B U$ for some $k \times n$ matrix U and some invertible symmetric $k \times k$ matrix B. Using the procedure outlined in the paragraphs following Lemma 2.9, we see that G is a blowup of a graph \tilde{G} corresponding to $\tilde{U}^t B \tilde{U}$, where \tilde{U} and B are defined as in the procedure. Lemma 2.9 then shows that $\tilde{G} \in \mathfrak{g}_k(\mathbb{F}_q)$.

Conversely, let G be a blowup of some graph in $\{H \cup K_1 \mid H \in \mathfrak{g}_k(\mathbb{F}_q)\}$ obtained by replacing each vertex v_i of H with a set of vertices V_i and K_1 with any number of vertices. Deleting isolated vertices of G does not change the minimum rank of G, so without loss of generality, we will assume that G has no isolated vertices (which implies that K_1 was replaced with an empty set of vertices). Let $\overline{x_1}, \overline{x_2}, \ldots, \overline{x_m}$ be the projective equivalence classes of $\mathbb{F}_q^k - \vec{0}$, with each x_i as a chosen representative from its class. Let $\tilde{U} = [x_1 \ x_2 \ \cdots \ x_m]$ and let B be an invertible symmetric $k \times k$ matrix such that $\tilde{U}^t B \tilde{U}$ corresponds to the graph H. Form the matrix $\hat{U} = [X_1 \ X_2 \ \cdots \ X_m]$ by replacing each column x_i of \tilde{U} with the block X_i , where the columns of X_i consist of $|V_i|$ copies of x_i . Then $\hat{U}^t B \hat{U}$ corresponds to G and rank $\hat{U}^t B \hat{U} \leq k$ since B has rank k. Thus $\operatorname{mr}(\mathbb{F}_q, G) \leq k$, so $G \in \mathcal{G}_k(\mathbb{F}_q)$. \Box

Now we will make this into a more explicit characterization of $\mathcal{G}_k(\mathbb{F}_q)$ by finding a suitable \mathcal{B} for any k and any q, thus enabling us to explicitly find $\mathfrak{g}_k(\mathbb{F}_q)$ for any k and any q.



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3. Congruence classes of symmetric matrices over finite fields. Symmetric matrices represent symmetric bilinear forms and play an important role in projective geometry. Two congruent symmetric matrices represent the same symmetric bilinear form with respect to different bases. Because of their fundamental importance, congruence classes of symmetric matrices over finite fields have been studied and characterized for a long time in projective geometry. In this section, we have distilled the pertinent proofs of these characterizations from [1], [23], and [16] to give a suitable \mathcal{B} for invertible symmetric $k \times k$ matrices over \mathbb{F}_q for any k and q. In the next section, we will expound more on the connection between the minimum rank problem and projective geometry.

We need the following elementary lemma.

LEMMA 3.1. If a symmetric matrix $B = \begin{bmatrix} C & D \\ D^t & E \end{bmatrix}$, where C is a square invertible matrix, then B is congruent to $\begin{bmatrix} C & O \\ O & E' \end{bmatrix}$, where O is a zero matrix and E' is a square symmetric matrix of the same order as E.

Proof. Let $R = C^{-1}D$ so that CR = D. Then

$$\begin{bmatrix} I & O \\ -R^t & I \end{bmatrix} \begin{bmatrix} C & D \\ D^t & E \end{bmatrix} \begin{bmatrix} I & -R \\ O & I \end{bmatrix} = \begin{bmatrix} C & D \\ -R^tC + D^t & -R^tD + E \end{bmatrix} \begin{bmatrix} I & -R \\ O & I \end{bmatrix}$$
$$= \begin{bmatrix} C & -CR + D \\ -R^tC + D^t & R^tCR - D^tR - R^tD + E \end{bmatrix}$$
$$= \begin{bmatrix} C & O \\ O & E - D^tR \end{bmatrix},$$

since $-CR + D = O = (-CR + D)^t = -R^tC + D^t$.

LEMMA 3.2. Every symmetric matrix over \mathbb{F}_q is congruent to a matrix of the form diag $(a_1, a_2, \ldots, a_s, b_1H_1, b_2H_2, \ldots, b_tH_t)$, where $a_i, b_i \in \mathbb{F}_q$, $H_i = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$, and s and t are nonnegative integers.

Proof. If B is the zero matrix, then the result is true.

If B is not the zero matrix, then the diagonal of B has a nonzero entry or there is some $a_{ij} \neq 0, i \neq j$, so that B has a principal submatrix of the form $\begin{bmatrix} 0 & a_{ij} \\ a_{ij} & 0 \end{bmatrix} = a_{ij}H$, where $H = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$.

In the first case, by using a suitable permutation, we may assume that $b_{11} \neq 0$. By Lemma 3.1, *B* is congruent to diag (b_{11}, B') .



In the second case, again by using a suitable permutation, we may assume that the upper left 2×2 principal submatrix is $a_{ij}H$. By Lemma 3.1, B is congruent to diag $(a_{ij}H, B')$.

Continue this process inductively with B'. Then, again using a suitable permutation, B is congruent to diag $(a_1, a_2, \ldots, a_s, b_1H, b_2H, \ldots, b_tH)$. \Box

We will now treat the even characteristic and odd characteristic cases separately.

3.1. Even characteristic. We first consider the case when \mathbb{F}_q has even characteristic. First, we need a well-known result.

LEMMA 3.3. Every element in a field of characteristic 2 is a square.

COROLLARY 3.4. Every symmetric matrix is congruent to diag $(I_s, H_1, H_2, \ldots, H_t)$.

Proof. By Lemma 3.2, a symmetric matrix A is congruent to a matrix

$$B = \text{diag}(a_1, a_2, \dots, a_s, b_1 H_1, b_2 H_2, \dots, b_t H_t).$$

Let

$$C = \operatorname{diag}(\frac{1}{\sqrt{a_1}}, \frac{1}{\sqrt{a_2}}, \dots, \frac{1}{\sqrt{a_s}}, \frac{1}{\sqrt{b_1}}I_2, \frac{1}{\sqrt{b_2}}I_2, \dots, \frac{1}{\sqrt{b_t}}I_2).$$

Then $C^t BC = \operatorname{diag}(I_s, H_1, H_2, \ldots, H_t)$. \Box

Let *B* be a symmetric matrix in \mathbb{F}_q . Then according to Corollary 3.4, *B* is congruent to a matrix $C = \text{diag}(I_s, H_1, H_2, \ldots, H_t)$, where each $H_i = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$. Either s = 0 or s > 0. If s > 0, then $\text{diag}(I_s, H_1, H_2, \ldots, H_t)$, and thus *B*, is congruent to I_k . To see this, let

$$A = \operatorname{diag}(1, H) = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix} \quad \text{and} \quad C = \begin{bmatrix} 1 & 1 & 1 \\ 1 & 0 & 1 \\ 0 & 1 & 1 \end{bmatrix}.$$

Then, since $\operatorname{char} \mathbb{F}_q = 2$,

$$C^{t}(AC) = \begin{bmatrix} 1 & 1 & 0 \\ 1 & 0 & 1 \\ 1 & 1 & 1 \end{bmatrix} \begin{bmatrix} 1 & 1 & 1 \\ 0 & 1 & 1 \\ 1 & 0 & 1 \end{bmatrix} = I_{3}.$$

If s = 0, then diag (H_1, H_2, \ldots, H_t) and B have even order and B is congruent to diag $(H_1, \ldots, H_{k/2})$.

The next lemma shows that these two cases are different.



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LEMMA 3.5. If a symmetric matrix B has a zero diagonal, then every matrix congruent to B has a zero diagonal.

Let B be a symmetric matrix having a zero diagonal. If v is the kth column of a matrix C, then the (k, k) entry of $C^t B C$ is $v^t B v$, which is zero, since

$$v^{t}Bv = \sum_{i,j} b_{ij}v_{i}v_{j} = \sum_{i} b_{ii}v_{i}^{2} + \sum_{i < j} b_{ij}(v_{i}v_{j} + v_{i}v_{j}) = \sum_{i} b_{ii}v_{i}^{2} = 0. \square$$

The results in this subsection give us the following lemma.

LEMMA 3.6. Let q be even. To determine $\mathfrak{g}_k(\mathbb{F}_q)$, we may take \mathcal{B} as follows: if k is odd, then $\mathcal{B} = \{I_k\}$; if k is even, then $\mathcal{B} = \{I_k, \operatorname{diag}(H_1, H_2, \ldots, H_{k/2})\}$, where $H_i = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$.

3.2. Odd characteristic. We now consider the case when \mathbb{F}_q has odd characteristic. We first need a well-known result.

LEMMA 3.7. If \mathbb{F}_q has odd characteristic and $\nu \in \mathbb{F}_q$, then there exists $c, d \in \mathbb{F}_q$ such that $c^2 + d^2 = \nu$.

Proof. Let $A = \{c^2 \mid c \in \mathbb{F}_q\}$ and $B = \{\nu - d^2 \mid d \in \mathbb{F}_q\}$. Since the map $\sigma: \mathbb{F}_q^{\times} \to \mathbb{F}_q^{\times}$ given by $\sigma: x \mapsto x^2$ has kernel $\{1, -1\}$, there are (q - 1)/2 squares in $\mathbb{F}_q \setminus \{0\}$. Including zero, there are then (q + 1)/2 squares in \mathbb{F}_q . Thus |A| = |B| = (q + 1)/2, so $A \cap B \neq \emptyset$, and $c^2 = \nu - d^2$ for some $c, d \in \mathbb{F}_q$. \Box

Since there are (q-1)/2 nonzero squares in \mathbb{F}_q , given a nonsquare $\nu \in \mathbb{F}_q$, the set $\{\nu b^2 \mid b \in \mathbb{F}_q, b \neq 0\}$ is a set of (q-1)/2 nonsquares in \mathbb{F}_q . Consequently, every nonsquare is equal to νb^2 for some $b \in \mathbb{F}_q$.

The matrix aH for any $a \in \mathbb{F}_q$ is congruent to a diagonal matrix:

[1	1] [0 0	$a] \begin{bmatrix} 1 \end{bmatrix}$	-1] _ [a	a	[1	-1	_	2a	0]	
$\lfloor -1$	1][a ($\begin{bmatrix} a \\ 0 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \end{bmatrix}$	1 .		a	-a	1	1		0	-2a	•

This fact combined with Lemma 3.2 shows that every symmetric matrix over \mathbb{F}_q is congruent to a diagonal matrix.

LEMMA 3.8. Every invertible symmetric $k \times k$ matrix B over \mathbb{F}_q is congruent to either I_k or diag (I_{k-1}, ν) , where ν is any nonsquare in \mathbb{F}_q .

Let C be an invertible diagonal matrix congruent to B, with $C = N^t B N$, and let ν be any nonsquare in \mathbb{F}_q .

By a permutation matrix P, let $D = P^t C P = \text{diag}(b_1^2, b_2^2, \dots, b_s^2, \nu c_1^2, \nu c_2^2, \dots, \nu c_t^2)$,



where the first s elements of the diagonal of D are squares in \mathbb{F}_q and the last t elements are nonsquares in \mathbb{F}_q .

Let
$$Q = \text{diag}(b_1^{-1}, b_2^{-1}, \dots, b_s^{-1}, c_1^{-1}, c_2^{-1}, \dots, c_t^{-1})$$
. Let $E = Q^t D Q = \text{diag}(I_s, \nu I_t)$.

Let $c, d \in \mathbb{F}_q$ such that $c^2 + d^2 = \nu$. Let

 $R = \nu^{-1} \left[\begin{array}{c} c & d \\ -d & c \end{array} \right].$

Since det $R = \nu^{-2}(c^2 + d^2) = \nu^{-1} \neq 0$, R is invertible. Note that

$$R^{t}(\nu I_{2})R = \nu R^{t}R = \nu \nu^{-2}(c^{2} + d^{2})I_{2} = I_{2}$$

If t is even, let $S = \text{diag}(I_s, R_1, R_2, \dots, R_{t/2})$, where $R_i = R$ for each i. Then $S^t ES = I_k$. If t is odd, let $S = \text{diag}(I_s, R_1, R_2, \dots, R_{(t-1)/2}, 1)$. Then $S^t ES = \text{diag}(I_{k-1}, \nu)$. \Box

The next lemma shows that these two cases are in fact different and gives a simple criteria to determine which congruence class any symmetric matrix is in.

LEMMA 3.9. If det B is a square (nonsquare) and \hat{B} is congruent to B, then det \hat{B} is a square (nonsquare).

Proof. Let $\hat{B} = C^t B C$. Then det $\hat{B} = (\det C)^2 (\det B)$. Thus det B is a square if and only if det \hat{B} is a square. \Box

Since det $I_k = 1$ is a square and det $(\text{diag}(I_{k-1}, \nu)) = \nu$ is a nonsquare, we can determine if a matrix is congruent to I_k or congruent to $\text{diag}(I_{k-1}, \nu)$ by whether the determinant is a square or not.

It appears then that $|\mathcal{B}| = 2$. However, we can do better in one case since we only are concerned with whether an entry of $U^t B U$ is zero or nonzero and not with the actual value of the entry.

DEFINITION 3.10. Let B and \hat{B} be matrices. If $\hat{B} = dC^tBC$ for some invertible matrix C and some nonzero constant d, then B and \hat{B} are projectively congruent.

Since multiplying by a nonzero constant preserves the zero/nonzero pattern in a matrix over a field, if B and \hat{B} are projectively congruent, then $U^t B U$ and $U^t \hat{B} U$ give isomorphic graphs.

LEMMA 3.11. If k is odd, then an invertible symmetric $k \times k$ matrix is projectively congruent to I_k .

Proof. Let $k = 2\ell - 1$. We can see that $\det(\nu \operatorname{diag}(I_{k-1}, \nu)) = \nu^{2\ell-1}\nu = \nu^{2\ell}$ is a square. Thus $\operatorname{diag}(I_{k-1}, \nu)$ is projectively congruent to I_k . \Box

The results in this subsection give us the following lemma.



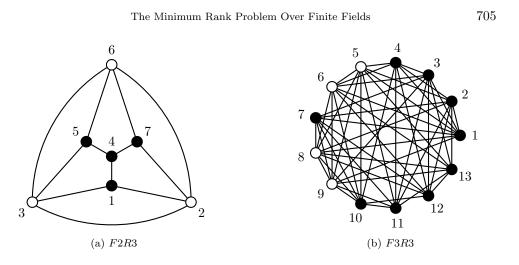


Fig. 3.1: Graphs in Corollaries 3.14 and 3.15

LEMMA 3.12. Let q be odd. To determine $\mathfrak{g}_k(\mathbb{F}_q)$, we may take \mathcal{B} as follows: if k is odd, then $\mathcal{B} = \{I_k\}$; if k is even, then $\mathcal{B} = \{I_k, \operatorname{diag}(I_{k-1}, \nu)\}$, where ν is any nonsquare in \mathbb{F}_q .

3.3. Summary. Combining Lemmas 3.6 and 3.12, the results of this section can be summarized as the following theorem.

THEOREM 3.13. The set $\mathfrak{g}_k(\mathbb{F}_q)$ is the set of graphs of the matrices in $\{U^t B U \mid B \in \mathcal{B}\}$, where the columns of U are a maximal set of nonzero vectors in \mathbb{F}_q^k such that no vector is a multiple of another and \mathcal{B} is given by:

- 1. if k is odd, $\mathcal{B} = \{I_k\}$. 2. if k is even and char $\mathbb{F}_q = 2$, $\mathcal{B} = \{I_k, \operatorname{diag}(H_1, H_2, \dots, H_{k/2})\}$, where $H_i = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$.
- 3. if k is even and char $\mathbb{F}_q \neq 2$, $\mathcal{B} = \{I_k, \operatorname{diag}(I_{k-1}, \nu)\}$, where ν is any non-square in \mathbb{F}_q .

3.4. Examples of characterizations. As special cases of Theorem 3.13, we present the following corollaries which calculate $\mathfrak{g}_k(\mathbb{F}_q)$ for several k and q. In the corollaries, we label a graph in $\mathfrak{g}_k(\mathbb{F}_q)$ using the pattern FqRk, signifying that it is a graph for the $\operatorname{mr}(\mathbb{F}_q, G) \leq k$ corollary. To compute these graphs, we used the software program Sage [32] and the Sage functions listed in Appendix A.

In these theorems, recall that K_1 does not have a loop.

COROLLARY 3.14. Let G be any simple graph. Let F2R3 be the graph in Fig-



ure 3.1(a). Then $mr(\mathbb{F}_2, G) \leq 3$ (i.e., $G \in \mathcal{G}_3(\mathbb{F}_2)$) if and only if G is a blowup graph of $F2R3 \cup K_1$.

As matrices over \mathbb{F}_2 , let

	0	1	0	1	0	1	1 -			[1	0	0 -	1
U =	0	0	1	1	1	1	0	and	B =	0	1	0	.
	1	1	1	1	0	0	0			0	0	1	

Then the graph F2R3 corresponds to the matrix

$$U^{t}BU = \begin{bmatrix} 1 & 1 & 1 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 & 1 & 1 & 0 \\ 0 & 1 & 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}. \square$$

Similarly, straightfoward matrix calculations give the following corollaries.

COROLLARY 3.15. Let G be any simple graph. Let F3R3 be the graph in Figure 3.1(b). Then $mr(\mathbb{F}_3, G) \leq 3$ (i.e., $G \in \mathcal{G}_3(\mathbb{F}_3)$) if and only if G is a blowup graph of $F3R3 \cup K_1$.

The next corollary gives the simplest previously-unknown result for which $\mathfrak{g}_k(\mathbb{F}_q)$ contains two graphs.

COROLLARY 3.16. Let G be any simple graph. Let F2R4A and F2R4B be the graphs in Figure 3.2. Then $mr(\mathbb{F}_2, G) \leq 4$ (i.e., $G \in \mathcal{G}_4(\mathbb{F}_2)$) if and only if G is a blowup graph of either $F2R4A \cup K_1$ or $F2R4B \cup K_1$.

4. Connection to projective geometry. As mentioned previously, the classifications of symmetric matrices in Section 3 are standard classification results in projective geometry. In this section, we first review appropriate terminology and highlight this connection to projective geometry. We will define slightly more terminology than is strictly necessary to help the reader see where these things fit into standard projective geometry. We then give some examples of how results in projective geometry can help us understand $\mathfrak{g}_k(\mathbb{F}_q)$ better. For further material, a definitive treatise on projective geometry is contained in the series [23] and [24].

4.1. Definitions and the connection. We start with basic definitions from projective geometry.



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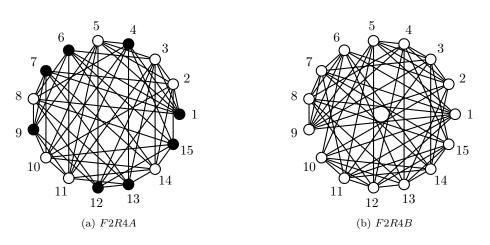


Fig. 3.2: Graphs in Corollary 3.16

DEFINITION 4.1. Let $V = \mathbb{F}_q^{n+1}$, the vector space of dimension n+1 over \mathbb{F}_q . For $x, y \in V - \vec{0}$, we define an equivalence relation by

$$x \sim y \iff x = cy$$
, where $c \in \mathbb{F}_q$ and $c \neq 0$.

Denote the equivalence class containing $x \in V - \vec{0}$ as $\bar{x} = \{cx \mid c \in \mathbb{F}_q \text{ and } c \neq 0\}$. Geometrically, we can think of the class \bar{x} as the set of non-origin points on a line passing through x and the origin in V. These equivalence classes form the *projective* geometry PG(n,q) of (projective) dimension n and order q. The equivalence classes are called the *points* of PG(n,q). Each subspace of dimension m+1 in V corresponds to a subspace of (projective) dimension m in PG(n,q). If a projective geometry has (projective) dimension 2, then it is called a *projective plane*.

Note that there is a shift by one in dimension between a vector space V and its subspaces and the projective geometry associated with V and its subspaces. To help the reader, we will use the nonstandard term *projective dimension* (or "pdim") when dealing with the dimension of a projective geometry.

DEFINITION 4.2. Let S be the set of subspaces of PG(n,q). A correlation $\sigma : S \to S$ is a bijective map such that for any subspaces $R, T \in S, R \subseteq T$ implies that $\sigma(T) \subseteq \sigma(R)$ and $\operatorname{pdim} \sigma(R) = n - 1 - \operatorname{pdim} R$. A polarity is a correlation σ of order 2 (i.e., $\sigma^2 = 1$, the identity map).

Note that any polarity σ maps points in S to hyperplanes (subspaces of projective dimension n-1 in S) and hyperplanes to points. Since $\sigma^2 = 1$, we have $Y = \sigma(\bar{x})$ if and only if $\sigma(Y) = \bar{x}$, so σ induces a bijection between points and hyperplanes. This bijection leads to the next definition.



DEFINITION 4.3. Let σ be a polarity on PG(n,q). Let \bar{x}, \bar{y} be points in PG(n,q). We say that $\sigma(\bar{x})$ is the *polar* (hyperplane) of \bar{x} and \bar{x} is the *pole* of $\sigma(\bar{x})$. If $\bar{y} \in \sigma(\bar{x})$, then $\bar{x} \in \sigma(\bar{y})$ and we say that \bar{x} and \bar{y} are *conjugate* points. If $\bar{x} \in \sigma(\bar{x})$, then we say that \bar{x} is *self-conjugate* or *absolute*. Similarly, if S is a subspace of PG(n,q), then S is *absolute* if $\sigma(S) \subseteq S$ or $S \subseteq \sigma(S)$. A subspace of PG(n,q) consisting of absolute points is called *isotropic*.

The next definition gives the connection with symmetric matrices.

DEFINITION 4.4. Let *B* be an $(n+1) \times (n+1)$ invertible symmetric matrix over \mathbb{F}_q . Define $\sigma : S \to S$ by $\sigma : R \mapsto R^{\perp}$, where the orthogonality relation is defined by the nondegenerate symmetric bilinear form represented by *B* (i.e., $R^{\perp} = \{\bar{y} \mid x^t By = 0 \text{ forall} \bar{x} \in R\}$). We call σ the polarity associated with *B*.

The fact that the σ in the previous definition is a polarity is easy to check.

Let M_1 and M_2 be symmetric matrices. Let σ_1 and σ_2 be the associated polarities, respectively. Two polarities are equivalent if the matrices are projectively congruent, i.e., σ_1 is equivalent to σ_2 if $M_1 = dC^t M_2 C$ for some nonzero d and invertible matrix C. Thus there is a unique polarity associated with each matrix given in Theorem 3.13.

We now summarize from [23, Section 2.1.5] the classification of polarities that are associated with symmetric matrices. Let B be an invertible symmetric matrix over \mathbb{F}_q . Let σ be the polarity associated with B.

• If q is odd, then σ is called an *ordinary polarity*.

If *B* has even order, then the associated polarity is either a *hyperbolic* polarity or an *elliptic* polarity. The correspondence between these types of polarities and the matrices in \mathcal{B} from Theorem 3.13(3) is slightly nontrivial and is summarized in [23, Corollary 5.19].

If B has odd order, then σ is a *parabolic* polarity, which corresponds to \mathcal{B} in Theorem 3.13(1).

- If q is even and $b_{ii} = 0$ for all i, then σ is a null polarity (or in alternate terminology, σ is a symplectic polarity). Note that this only occurs when B has even order since otherwise B is not invertible. This case corresponds to the non-identity matrix in the \mathcal{B} in Theorem 3.13(2).
- If q is even and there is some $b_{ii} \neq 0$, then σ is a *pseudo-polarity*. This case corresponds to the identity matrix in \mathcal{B} in Theorem 3.13(1) or (2).

We pause to note that there are polarities that are not associated with symmetric matrices. However, since we are only concerned about symmetric matrices, we will restrict ourselves to this case. Information about polarities not associated with symmetric matrices may also be found in [23].



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We now examine the connection to graphs by recalling the definition of a polarity graph.

DEFINITION 4.5. Let *B* be an invertible symmetric $(n + 1) \times (n + 1)$ matrix over \mathbb{F}_q and let σ be the associated polarity. The *polarity graph* of σ has as its vertices the points of PG(n,q) and as its edges $\{\bar{x}\bar{y} \mid x^tBy = 0\}$. In a polarity graph, \bar{x} is adjacent to \bar{y} exactly when \bar{x} and \bar{y} are conjugate (i.e., x and y are orthogonal with respect to *B*). In standard literature, loops are not allowed in polarity graphs. However, for our purposes, loops convey needed information, so a vertex \bar{x} in a polarity graph has a loop if and only if \bar{x} is absolute (i.e., $x^tBx = 0$, where *B* is an invertible symmetric matrix associated with the polarity).

In Theorem 3.13, the vertices of a graph in $\mathfrak{g}_k(\mathbb{F}_q)$ represent the points of the projective geometry PG(k-1,q) and an edge is drawn if the corresponding points are *not* conjugate (i.e., $x^t By \neq 0$). Thus, the graphs in Theorem 3.13 are exactly the complements of polarity graphs. Recall that when dealing with looped graphs, a vertex is looped in the complement of a graph if and only if it is nonlooped in the original graph.

Using this connection, we can restate Theorem 3.13:

THEOREM 4.6. The set $\mathfrak{g}_k(\mathbb{F}_q)$ is the set of complements of the (looped) polarity graphs of the polarities on PG(k-1,q) that are associated with symmetric matrices.

4.2. Consequences of the connection. With the main theorem stated as in Theorem 4.6, we can use a variety of known results about polarity graphs to derive results about graphs in $\mathfrak{g}_k(\mathbb{F}_q)$. In this section, we list a few consequences of Theorem 4.6.

An elementary result in projective geometry gives us the size of the graphs in $\mathfrak{g}_k(\mathbb{F}_q)$. While this result could have been realized from the statement in Theorem 3.13, it also naturally follows as a consequence of Theorem 4.6.

THEOREM 4.7. Every graph in $\mathfrak{g}_k(\mathbb{F}_q)$ has $\frac{q^k-1}{q-1}$ vertices.

Proof. There are $q^k - 1$ vectors in $\mathbb{F}_q^k - \vec{0}$. Since there are q - 1 nonzero constants in \mathbb{F}_q , there are q - 1 elements in each equivalence class in PG(k - 1, q), so there are $\frac{q^k - 1}{q - 1}$ points in PG(k - 1, q). \Box

The following observation follows directly from Theorem 4.6 and restates the criteria for an edge in a graph in $\mathfrak{g}_k(\mathbb{F}_q)$ in several ways.

OBSERVATION 4.8. Let $G \in \mathfrak{g}_k(\mathbb{F}_q)$ and let u and v be (not necessarily distinct) vertices in G. Let σ be the polarity corresponding to G and let B be an invertible symmetric matrix corresponding to σ . Then uv is an edge in G if and only if:



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1. $u^t Bv \neq 0$ (equivalently, $v^t Bu \neq 0$), or equivalently,

2. u and v are not conjugate points, or equivalently,

3. $u \notin \sigma(v)$ (equivalently, $v \notin \sigma(u)$).

COROLLARY 4.9. A graph $G \in \mathfrak{g}_k(\mathbb{F}_q)$ is regular of degree q^{k-1} (using the convention that a loop adds one to the degree of a vertex). Let $v \in G$ and let σ be the polarity associated with G. Since the hyperplane $\sigma(v)$ contains $\frac{q^{k-1}-1}{q-1}$ points, this is the degree of a v in the complement of G. Thus the degree of v in G is

$$\frac{q^k - 1}{q - 1} - \frac{q^{k - 1} - 1}{q - 1} = q^{k - 1}. \qquad \Box$$

In light of Observation 4.8, determining the numbers of looped and nonlooped vertices in G is equivalent to finding the numbers of absolute points of the polarities of PG(k-1,q).

THEOREM 4.10. Let \mathbb{F}_q be a finite field having characteristic 2. One graph in $\mathfrak{g}_k(\mathbb{F}_q)$ will have $\frac{q^{k-1}-1}{q-1}$ nonlooped vertices. If k is even, then the additional graph in $\mathfrak{g}_k(\mathbb{F}_q)$ will have all nonlooped vertices.

Proof. In a field of characteristic 2, since

$$x^{t}Bx = \sum_{i,j} b_{ij}x_{i}x_{j} = \sum_{i} b_{ii}x_{i}^{2} + \sum_{i < j} b_{ij}(x_{i}x_{j} + x_{i}x_{j}) = \sum_{i} b_{ii}x_{i}^{2} = \left(\sum_{i} \sqrt{b_{ii}}x_{i}\right)^{2},$$

a point \bar{x} is absolute if and only if $\sum_i \sqrt{b_{ii}} x_i = 0$.

In a pseudo-polarity, the set of absolute points is the hyperplane $\sum_i \sqrt{b_{ii}}x_i = 0$. Since a hyperplane of PG(k-1,q) is a projective geometry of projective dimension k-2, there are $\frac{q^{k-1}-1}{q-1}$ nonlooped vertices in this graph.

In a null polarity, $b_{ii} = 0$ for all i. Therefore every vertex is nonlooped (i.e., there are $\frac{q^k - 1}{a - 1}$ nonlooped vertices). A null polarity occurs when k is even. \Box

For the odd characteristic case, we will directly apply a standard result in projective geometry about the number of absolute points in ordinary polarities.

THEOREM 4.11 ([24, Theorem 22.5.1(b)]). Let q be odd. Then the number of absolute points in a polarity in PG(k-1,q) is given by:

• $\frac{(q^m-1)(q^{m-1}+1)}{q-1}$ or $\frac{(q^m+1)(q^{m-1}-1)}{q-1}$ if k = 2m is even; • $\frac{q^{2m}-1}{q-1}$ if k = 2m+1 is odd.

COROLLARY 4.12. Let q be odd. If k = 2m is even, then the two graphs in $\mathfrak{g}_k(\mathbb{F}_q)$ will have $\frac{(q^m-1)(q^{m-1}+1)}{q-1}$ and $\frac{(q^m+1)(q^{m-1}-1)}{q-1}$ nonlooped vertices, respectively. If k = 2m + 1 is odd, then the graph in $\mathfrak{g}_k(\mathbb{F}_q)$ will have $\frac{q^{2m}-1}{q-1}$ nonlooped vertices.



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We conclude by applying a few standard results for polarities over PG(2,q) (a projective plane) to give results about $\mathfrak{g}_3(\mathbb{F}_q)$ and the minimum rank problem. We note that the polarity graphs of PG(2,q) for any q are the Erdős-Rényi graphs from extremal graph theory (see [19], [20], or [12]). For a survey of interesting properties of the Erdős-Rényi graphs and their subgraphs, see [30] or [34, Chapter 3].

THEOREM 4.13. If $G \in \mathfrak{g}_3(\mathbb{F}_q)$, then the nonlooped vertices in G form a clique.

Proof. Suppose that u and v are distinct nonadjacent nonlooped vertices in G. Then u and v are absolute vertices and $u \in \sigma(u) \cap \sigma(v)$ and $v \in \sigma(u) \cap \sigma(v)$. This is a contradiction since the intersection of any two distinct lines in PG(2,q) is a single point. \Box

If $G \in \mathfrak{g}_3(\mathbb{F}_q)$, the formulas in Theorem 4.10 and Corollary 4.12 imply that G has q + 1 nonlooped vertices. This combined with Corollary 4.9 and Theorem 4.13 gives the following corollary.

COROLLARY 4.14. If $G \in \mathfrak{g}_3(\mathbb{F}_q)$, then each nonlooped vertex is adjacent to q nonlooped vertices and $q^2 - q$ looped vertices.

Theorem 4.13 also gives us the following theorem.

THEOREM 4.15. Let $G = K_{s_1,s_2,...,s_n}$, a simple complete multipartite graph. If $q \ge n-1$, then $mr(\mathbb{F}_q, G) \le 3$.

Proof. Let $G = K_{s_1,s_2,...,s_n}$. Then G is a blowup graph of K_n , where each vertex of K_n is nonlooped. Since the graph in $\mathfrak{g}_3(\mathbb{F}_q)$ contains a clique of q+1 nonlooped vertices, if $q+1 \geq n$, then G is a blowup graph of the graph in $\mathfrak{g}_3(\mathbb{F}_q)$. \Box

We can now construct an interesting family of simple graphs.

THEOREM 4.16. For every integer $n \ge 1$, let G_n be a simple complete multipartite graph $H_1 \lor H_2 \lor \cdots \lor H_n$ where each H_i is an independent set with $s_i > (n-1)^2$ vertices. We then have $mr(\mathbb{F}_q, G_n) \le 3$ if and only if $q \ge n-1$.

Proof. If $q \ge n-1$, then $mr(\mathbb{F}_q, G_n) \le 3$ by Theorem 4.15.

Conversely, let q < n - 1. Let I be the graph in $\mathfrak{g}_3(\mathbb{F}_q)$ and let I_1 and I_2 be the subgraphs of I induced by the looped and nonlooped vertices of I, respectively. Since I_1 has q^2 vertices, any blowup of I_1 containing more than q^2 vertices will contain an edge by the pigeon-hole principle. Since the vertices in each H_i form an independent set of size $s_i > (n-1)^2 > q^2$, at least one vertex in each H_i must be a blowup of a vertex in I_2 . Furthermore, since the vertices of each H_i have the same neighbors, we can assume without loss of generality that all of the vertices of each H_i are blowups of vertices of I_1 . Thus G_n is a blowup of I_2 . However, any blowup of I_2 will be of the form $K_{t_1,t_2,...,t_{q+1}}$ since I has q+1 nonlooped vertices, but G_n is not of this form



since q + 1 < n. \Box

5. Conclusion. We have succeeded in classifying the structure of graphs in $\mathcal{G}_k(\mathbb{F}_q)$ for any k and any q. We have also shown how this classification relates to projective geometry. We have applied a few results of projective geometry to give results in the minimum rank problem.

We conclude with a short list of open questions and topics for further investigation. First, there are many results about polarity graphs that could potentially yield results for the minimum rank problem. What other facts from projective geometry can be applied to give results in the minimum rank problem over finite fields?

The structural characterization in this paper gives rise to a theoretical procedure for determining the minimum rank of any graph over a finite field. How can this procedure be efficiently implemented? How can the results of Ding and Kotlov [18] be combined with the classification in this paper to yield results on minimal forbidden subgraphs describing $\mathcal{G}_k(\mathbb{F}_q)$? The author has implemented such an algorithm and has some preliminary results on the numbers of forbidden subgraphs describing $\mathcal{G}_k(\mathbb{F}_q)$ for different values of k and q.

Finally, there is still ongoing research investigating the structure of polarity graphs. For example, Jason Williford [34], Michael Newman, and Chris Godsil [21] have recently investigated the sizes of independent sets in polarity graphs. Are there results in the minimum rank problem that would aid in answering questions about the structure of polarity graphs?

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Appendix A. Sage code to generate graphs.

```
# This code is written for Sage 2.8.15.
# See http://www.sagemath.org/
```

```
def bilinearforms(F,mr):
        # Construct a matrix space for our bilinear forms
5
        MSpace = MatrixSpace(F,mr)
        # The identity matrix is always
        \# a congruence class representative
        forms = [MSpace.identity_matrix()]
        \# Add the extra matrices in the even rank cases
10
        if(mod(mr,2)==0): \# even rank
             if (F. characteristic ()==2): \# characteristic 2
                 \# Add \operatorname{diag}(H_1, H_2, \ldots, H_{\mathrm{mr}/2})
                 hyperbolic = matrix (F, [[0, 1], [1, 0]])
                 hyperbolic_form = matrix(F,[])
15
                 for i in (1.. \operatorname{Integer}(\operatorname{mr}/2)):
                      hyperbolic_form = hyperbolic_form \setminus
                                             . block_sum(hyperbolic)
                 forms.append(hyperbolic_form)
             else: # odd characteristic
20
                 # Add diag(I_{n-1}, \nu), where \nu is a non-square
                 nonidentity_form = MSpace.identity_matrix()
                 # Find a non-square
                 for nu in F:
                      if not nu.is_square():
^{25}
```



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break if nu.is_square(): **raise** NotImplementedError, \ "Cannot_find_a_non-square_in_field." nonidentity_form [mr-1, mr-1] = nu30 forms.append(nonidentity_form) return forms **def** get_matrices (F,mr): # U has one vector for every equivalence class 35 # in PG(mr-1,q)U = matrix(F, [list(v)])for v in ProjectiveSpace(mr-1,F)]) .transpose() B = bilinear forms(F,mr)40 return U, B **def** $get_graphs(F,mr)$: $U, B = get_matrices(F, mr)$ product_matrices = [U.transpose()*b*U for b in B] 45graphs = [Graph(m) for m in product_matrices] **for** i **in** range(len(graphs)): graphs [i].loops(true); graphs[i].add_edges([[j,j] for j **in** range(len(graphs[i])) \ 50if product_matrices[i][j,j] != 0])return graphs **def** show_graphs(F,mr): 55for g in get_graphs(F,mr): # Vertices with loops are black, others are white $vcolors = \{ 'black': g.loop_vertices(), \setminus \}$ 'white': [i for i in g.vertices() if i not in g.loop_vertices()]} 60 g.show(layout='circular', vertex_colors=vcolors, \setminus vertex_labels=false)

To retrieve the matrices for graphs in $g_3(\mathbb{F}_2)$: ⁶⁵ U, B = get_matrices (F=FiniteField (2), mr=3)



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To retrieve the graphs in $g_3(\mathbb{F}_2)$: graph_list = get_graphs(F=FiniteField(2), mr=3)

70 # To display the graphs with vertices colored appropriately: show_graphs (F=FiniteField (2), mr=3)