

ON MINIMAL RANK OVER FINITE FIELDS*

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Abstract. Let F be a field. Given a simple graph G on n vertices, its *minimal rank* (with respect to F) is the minimum rank of a symmetric $n \times n$ F -valued matrix whose off-diagonal zeroes are the same as in the adjacency matrix of G . If F is finite, then for every k , it is shown that the set of graphs of minimal rank at most k is characterized by finitely many forbidden induced subgraphs, each on at most $(\frac{|F|^k}{2} + 1)^2$ vertices. These findings also hold in a more general context.

Key words. Minimal rank, Forbidden induced subgraph, Critical graph.

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1. Introduction. Let F be a field. Given a simple graph G on n vertices, we say that a symmetric $n \times n$ F -valued matrix *represents* G if its off-diagonal zeroes are the same as in the adjacency matrix of G . The *min rank* [w.r.t. F] of G is the minimum rank [over F] of the matrices representing G . The reader is referred to [1] for motivation and importance of min rank.

For a fixed integer k , denote $\mathfrak{G}_k$ the class of all graphs of min rank at most k . It is obvious that $\mathfrak{G}_k$ is closed under vertex deletion. Call a graph *k-critical* if it is minimal (w.r.t. vertex deletion) of rank larger than k . Clearly, $\mathfrak{G}_k$ is characterized by the k -critical graphs as forbidden induced subgraphs.

It is a celebrated result of Robertson and Seymour [3] that any class of graphs closed under taking *minors* is characterized by a *finite* set of forbidden minors. At the 2005 Oberwolfach graph theory workshop, Hein van der Holst asked if there are only finitely many k -critical graphs for any k and F . (Barrett, van der Holst, and Loewy [2, 1] had recently confirmed this for $k \leq 2$ and any F .) When F is finite, we answer this question affirmatively, by providing an upper bound on the size of a k -critical graph. In fact, our arguments hold in a more general context. In the next section, we define k -critical graphs w.r.t. an arbitrary collection $\mathfrak{M}_k$ of matrices of rank at most k and show that the number of such graphs is finite as long as there is $c \in \mathbb{N}$ such that each matrix in $\mathfrak{M}_k$ uses at most c distinct elements from its field. If, in addition, $\mathfrak{M}_k$ is closed under row-and-column-duplication, the number of k -critical graphs can be bounded in terms of c and k . We derive such bounds in Section 3.

2. Definitions and Main Result. For a fixed $k \in \mathbb{N} \cup \{0\}$, let $\mathfrak{S}_k$ denote the set of all pairs (M, F) where F ranges over all fields and M is an F -valued symmetric matrix with $\text{rk}_F(M) \leq k$. Fix some non-empty $\mathfrak{M}_k \subseteq \mathfrak{S}_k$. Denote by $\mathfrak{G}_k$ the set of graphs represented by matrices in $\mathfrak{M}_k$. As before, call a graph G *k-critical* if

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$G \notin \mathfrak{G}_k$ while every proper induced subgraph of G is in $\mathfrak{G}_k$. It is easy to see that $\mathfrak{G}_k$ is characterized by $[k\text{-critical}]$ graphs as forbidden induced subgraphs iff it is closed under vertex deletion.

For a matrix A , denote by $c(A)$ the number of distinct entries in A , and for any non-empty $\mathfrak{A} \subseteq \mathfrak{G}_k$, set $c(\mathfrak{A}) := \sup\{c(A) : (A, F) \in \mathfrak{A}\}$.

OBSERVATION 2.1. *A matrix A of rank r has at most $c(A)^r$ distinct rows.*

Proof. W.l.o.g. assume that the first r columns of A are linearly independent. Then two rows of A are identical iff they agree on the first r coordinates. $\square$

Given disjoint graphs G and H , by replacing a vertex v of G with H we mean, as is common, the disjoint union of $G - v$ and H plus the edges $\{uv : uv \in G, v \in H\}$. If a matrix M represents G and the last $t \geq 2$ rows of M are identical then G can be obtained from the graph $G - \{n - t + 2, \dots, n\}$ by replacing vertex $n - t + 1$ with either the clique K_t (if $M_{nn} \neq 0$) or the independent set $\bar{K}_t$ (if $M_{nn} = 0$). For $m \in \mathbb{N}$, call a graph m -sprawling if it can be obtained from a graph on at most m vertices by replacing each vertex with either a clique or an independent set.

COROLLARY 2.2. *Let $c := c(\mathfrak{M}_k) < \infty$. Then every graph in $\mathfrak{G}_k$ is c^k -sprawling. Consequently, every k -critical graph is $(c^k + 1)$ -sprawling. $\square$*

LEMMA 2.3. *Given $m \in \mathbb{N}$, an infinite sequence of distinct m -sprawling graphs has an infinite subsequence ascending w.r.t. induced-subgraph inclusion.¹*

Proof. Fix an infinite sequence, $\mathbf{s}$, of $[distinct]$ m -sprawling graphs. As there are finitely many graphs on at most m vertices, $\mathbf{s}$ has an infinite subsequence, $\mathbf{s}'$, of graphs which can be sprawled from the same graph, G , on some $n \leq m$ vertices. Further, as there are finitely many (namely, 2^n) choices of whether to use a clique or an independent set at each vertex of G , $\mathbf{s}'$ contains an infinite subsequence $\mathbf{s}''$ of graphs H_i ($i \in \mathbb{N}$) for which such choices coincide. Now, each H_i can be described by a string of n natural numbers $(a_{i1}, \dots, a_{in})$ where a_{ij} is the number of vertices by which vertex j of G was replaced in obtaining H_i . Notice that an infinite sequence of natural numbers contains a monotone non-decreasing infinite subsequence. By sequentially applying this argument n times, we find that $\mathbf{s}''$ has an infinite subsequence, $\{H_{i_t} : t \in \mathbb{N}\}$, such that, for each $j = 1, \dots, n$, the sequence $\{a_{i_t j}\}$ is monotone non-decreasing. But then, the sequence $\{H_{i_t}\}$ itself is monotone non-decreasing under induced-subgraph inclusion. $\square$

COROLLARY 2.4. *If $c := c(\mathfrak{M}_k) < \infty$, then there are finitely many k -critical graphs.*

Proof. An infinite sequence of distinct k -critical (and thus, $(c^k + 1)$ -sprawling) graphs would contain an ascending $[infinite]$ subsequence, contrary to the definition of k -criticality. $\square$

¹A partially ordered set whose every infinite sequence of distinct elements has an increasing subsequence is called *well-partially-ordered*, cf. e.g. [4]. In this terminology, the lemma becomes: *the set of m -sprawling graphs is well-partially-ordered w.r.t. induced-subgraph inclusion.*

We summarize our findings as follows.

THEOREM 2.5. *If $c(\mathfrak{M}_k) < \infty$ and $\mathfrak{G}_k$ is closed under vertex deletion, then $\mathfrak{G}_k$ is characterized by finitely many k -critical graphs as forbidden induced subgraphs.*

3. Upper Bound. In the previous section, we saw that once $c(\mathfrak{M}_k) < \infty$, the number of k -critical graphs is finite. However, in the absence of other requirements, this number need not be bounded by any function of c and k . Indeed, fix some $N \in \mathbb{N}$ and let $\mathfrak{M}_1$ consist of all the rank-at-most-1 symmetric 0-1 matrices of size less than $N \times N$. Then the N graphs $\{K_a \dot{\cup} \overline{K}_{N-a} : a = 1, \dots, N\}$ (the disjoint unions of the a -clique and $N - a$ independent vertices) are 1-critical.

In this section, we bound the number of k -critical graphs in terms of c and k under the following additional constraint on $\mathfrak{M}_k$:

- (*) *If $(M, F) \in \mathfrak{M}_k$ then also $(M', F) \in \mathfrak{M}_k$ where M' has two identical rows i and j and M can be obtained by deleting row j and column j in M' .*

THEOREM 3.1. *If (*) holds and $c := c(\mathfrak{M}_k) < \infty$, then $|G| \leq (\frac{c^k}{2} + 1)^2$ for a k -critical G .*

We precede the proof of Theorem 3.1 by the following discussion. Two vertices of a graph G are *twins* if they have the same neighbors in the rest of G (the two vertices themselves can be either adjacent or independent). It is easy to see that a twin relationship is transitive: if u and v are twins, and v and w are twins, then u and w are also twins and, moreover, the triple $\{u, v, w\}$ spans either a clique or an independent set. (We will call such a triple of pairwise-twin vertices a *triplet*.) Thus, the twin relationship induces a partition, $T(G)$, of the vertex set of G into *twin classes*.

Fix a vertex $v \in G$ and consider a partition $T' := T(G) - v$ of the vertex set of $G - v$. It is easy to see that: (a) T' is a refinement of $T(G - v)$; (b) more specifically, every set in $T(G - v)$ is a [disjoint] union of one or two sets in T' ; (c) in particular, if G has no twins then $G - v$ has no triplet; (d) if v has exactly one twin, v' , in G then T' and $T(G - v)$ coincide except perhaps on the twin class of v' in $G - v$; (e) if v has at least two twins in G then $T' = T(G - v)$. In addition, (f) if G has no triplet then v can be chosen so that $G - v$ has no triplet either. Indeed, if G is twin-free then (f) is trivially true by (c). Thus, assume G is not, and let v and v' be twins in G . Then, by (d), the only triplet $G - v$ may have is of the form $\{v', u, u'\}$, where u and u' are twins in G . Further, the subgraph induced by v, v', u, u' has either no edge except vv' or all edges except vv' , and v, v', u, u' have the same neighbors in the rest of G . But then, $T(G) - u = T(G - u)$, as required.

Proof of Theorem 3.1: Clearly, we may assume that $\text{rk}_F(M) = k$ for some $(M, F) \in \mathfrak{M}_k$. Suppose first that $k \geq 2$. Then necessarily $c \geq 2$. If G has no triplet then, by (f) and Observation 2.1, $n - 1 \leq 2c^k \leq (1 + \frac{c^k}{4})c^k = (\frac{c^k}{2} + 1)^2 - 1$, as required. Hence, we can assume w.l.o.g. that the last $t \geq 3$ vertices of G form its largest twin class, C . By criticality of G , there is $(M, F) \in \mathfrak{M}_k$ such that M represents $G - n$. Consider the diagonal entry M_{ii} of M for some $i > n - t$. If C is an independent set in G and $M_{ii} = 0$ or if C is a clique and $M_{ii} \neq 0$ then the symmetric $n \times n$ matrix M'

whose rows i and n are identical and whose upper-left $(n-1) \times (n-1)$ corner is M , represents G and, by $(*)$, $(M', F) \in \mathfrak{M}_k$. This is a contradiction. Denoting by D the lower-right $(t-1) \times (t-1)$ corner of M , we thus have:

- (†) *either D has no off-diagonal zero and its diagonal is all-zero, or D is all-zero except its diagonal has no zero.*

In particular, no two among the last $t-1$ rows of M are identical. Hence, by Observation 2.1, among the first $n-t$ rows of M , at most $c^k - (t-1)$ are distinct, whence at least $\frac{n-t}{c^k - (t-1)}$ are identical. The latter is a lower bound on the size of a twin class in $G - n$ and thus, by (e), also on t . We have $\frac{n-t}{c^k - (t-1)} \leq t$ whence $n \leq t(c^k + 2 - t) \leq (\frac{c^k + 2}{2})^2$, as required.

It remains to consider the case $k \leq 1$. If $k = 0$ then $G = K_2$, and if $k = 1$ and $c = 1$ then $G = \overline{K}_2$, both in accord with what claimed. Thus assume that $k = 1$ and $c \geq 2$ but, contrary to the claim, $|G| \geq 5$. Notice that G contains neither the path P_3 nor $K_2 \dot{\cup} K_2$ as an induced subgraph as neither is in $\mathfrak{G}_1$. Hence, the non-isolated vertices of G , if any, form a clique. But then, in the above notation, $t \geq 3$ and, by (†), the lower-right 2×2 corner of M is of F -rank 2. This is a contradiction. $\square$

We can improve the bound of Theorem 3.1 by a more diligent accounting. For the sake of brevity, we only establish the asymptotic version of such an improvement.

THEOREM 3.2. *Under the assumptions of Theorem 3.1, $|G| = O(2^k + (c-1)^{2k})$.*

Proof. We utilize the notation of the proof of Theorem 3.1. As remarked in that proof, $n = O(c^k)$ if $t \leq 2$, in accord with our claim. Thus, assume $t \geq 3$.

Set $r := \text{rk}_F(D)$ and let, w.l.o.g., the last r columns of D be linearly independent. As in Observation 2.1, we can bound the number of distinct rows in M by considering its last r columns. Namely, among the first $n-t$ rows of M , the distinct ones come from at most c^{k-r} distinct rows whose last r components are all-zero and at most $c^{k-r}(c-1)^r$ distinct rows whose last r components have no zero. Proceeding as in the proof of Theorem 3.1, we find that $n \leq t(c^{k-r}((c-1)^r + 1) + 1)$. Further, the number $t-1$ of [distinct] rows of D is, again by (†), at most $(c-1)^r + r$, whence

$$(\dagger) \quad n \leq ((c-1)^r + r + 1)(c^{k-r}((c-1)^r + 1) + 1).$$

If $c \geq 3$ then (†) implies $n = O(c^{k-r}(c-1)^{2r})$ which, in turn, implies $n = O((c-1)^{2k})$, and if $c = 2$, (†) becomes $n \leq (r+2)(2^{k-r+1} + 1) = O(2^k)$ —all in accord with the claim. $\square$

4. Concluding Remarks. The following improvements can be made to the argument for Theorem 3.2 should one desire to derive an exact (rather than asymptotic) bound there. Below, we assume that $(*)$ holds and G is a k -critical graph.

1. If a twin-class C of G is an independent set then, by (†), $|C| \leq k+1$. Moreover, $|C| = k+1$ may hold only if the non-neighbors of C form an independent set themselves: this is because the direct sum $A \oplus B$ of matrices A and B satisfies $\text{rk}(A \oplus B) = \text{rk}(A) + \text{rk}(B)$. In particular, G has at most $k-1$ isolated vertices unless $G = \overline{K}_k \dot{\cup} K_m$ for some $m \in \mathbb{N}$.

2. Similarly, if a twin-class C of G is a clique and $c = 2$ then, in the notation of the proof of Theorem 3.2, $r \leq |C| - 1 \leq r + (c - 1)^r = r + 1$ and it is easy to see that $|C| - 1 = r + 1$ (equivalently, that the first row of D is a linear combination of the rest) iff $\text{char}(F)$ divides r ; further, $r \leq k$ and the equality may hold only if the non-neighbors of C form an independent set.
3. More generally, let $f(r, c)$ denote an upper bound on the size of an F -valued symmetric matrix D described by $(\dagger)$ such that $c(D) \leq c$ and $\text{rk}_F(D) = r$. In the proof of Theorem 3.2, we implicitly set $f(r, c) := (c - 1)^r + r$ and concluded that $n = O(f(r, c)c^{k-r}(c - 1)^r)$. Thus, a tighter f , should one exist, would lead to a tighter bound in Theorem 3.2. (In the previous remark then, we showed how to improve $f(r, 2)$ in particular cases.)
4. Finally, in the notation of the proof of Theorem 3.2, the upper bound on the number of distinct rows among the first $n - t$ rows of M can be reduced by $t - 1 - r$, the number of those distinct rows with no zero (or no non-zero) in the last r positions which are already used among the last $t - 1$ rows of M .

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