

RECOGNITION OF HIDDEN POSITIVE ROW DIAGONALLY DOMINANT MATRICES*

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Abstract. A hidden positive row diagonally dominant (hprdd) matrix is a square matrix A for which there exist square matrices C and B so that $AC = B$ and each diagonal entry of B and C is greater than the sum of the absolute values of the off-diagonal entries in its row. A linear program with $5n^2 - 4n$ variables and $2n^2$ constraints is defined that takes as input an $n \times n$ matrix A and produces C and B satisfying the above conditions if and only if they exist. A 4×4 symmetric positive definite matrix that is not an hprdd matrix is presented.

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1. Introduction and Terminology. An $n \times n$ matrix $A = (a_{ij})$ with real entries is called *row diagonally dominant* if, for $i = 1, \dots, n$, we have

$$|a_{ii}| > \sum_{i \neq j} |a_{ij}|.$$

The matrix A is called a *P-matrix* if the principal minors of A are all positive. The problem of determining if A is not a *P-matrix* is known to be *NP*-complete [2]. Two of the most important subclasses of the class of *P-matrices* are the class of positive definite matrices and the class of *M-matrices*. *M-matrices* are *P-matrices* that have non-positive off-diagonal elements. Membership in each of these subclasses can be checked in polynomial time. The class of *M-matrices* has been generalized to the class of *hidden M* matrices. A *P-matrix* A is *hidden M* if there are matrices B and C such that $AC = B$, where C and B both have non-positive off-diagonal entries and C is an *M-matrix*. Membership in the *hidden M* class can also be tested in polynomial time (see [4]). It has recently been (see [7]) proved that A is a *P-matrix* whenever one can find row diagonally dominant matrices B and C , with positive diagonal entries, so that $AC = B$. We will call matrices satisfying this sufficient condition for *P-matrices* *hidden positive row diagonally dominant*, or *hprdd*.

Tsatsomeros [7] proved that every *hidden-M* matrix is *hprdd*. No example of a *P-matrix* that is not *hprdd* was known until now, and there was some speculation (see [7]) that there might not be such a matrix. Part of the difficulty in producing such a matrix was that no efficient algorithm had been published for checking the *hprdd* property. We present a linear program with $5n^2 - 4n$ variables and $2n^2$ constraints that produces the required B and C if they exist. If the input matrix is not *hprdd*, the solution to the dual linear program provides a proof that the matrix is not *hprdd*. Experimentation

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with randomly generated positive definite matrices yielded the following example of a positive definite matrix that is not *hprdd*:

$$\begin{bmatrix} 4 & 4 & 7 & -9 \\ 4 & 16 & -7 & -2 \\ 7 & -7 & 30 & -24 \\ -9 & -2 & -24 & 25 \end{bmatrix}.$$

2. The Linear Program. The set of vectors K in $\mathbb{R}^n$ satisfying an inequality $x_i > \sum_{i \neq j} |x_j|$ is the interior of a cone that has $2(n-1)$ extreme rays, each of the form $e_i \pm e_j$ for standard basis vectors $e_i \neq e_j$. The set of pairs (C, B) of $n \times n$ matrices can be identified with $\mathbb{R}^{2n^2}$. Define $U = \{(C, B) \in \mathbb{R}^{2n^2} : B \text{ and } C \text{ are both positive row diagonally dominant}\}$. Then U is the Cartesian product of $2n$ copies of K . It is the interior of a convex cone that has the $4n(n-1)$ extreme rays of the form $(E_{ii} \pm E_{ij}, 0)$, $i \neq j$ or $(0, E_{ii} \pm E_{ij})$, $i \neq j$ where E_{ij} is the $n \times n$ matrix that has 1 in entry (i, j) and 0 everywhere else.

Let A be an $n \times n$ matrix. Define $V = \{(C, B) \in \mathbb{R}^{2n^2} : AC = B\}$. Note that V is a subspace of $\mathbb{R}^{2n^2}$ that is spanned by the n^2 vectors of the form (E_{ij}, AE_{ij}) . A matrix A is *hprdd* if and only if $V \cap U$ is not empty.

In order to define a linear program to check for the *hprdd* property of a matrix A , let H be a $2n^2 \times n^2$ matrix that has as its columns the vectors (E_{ij}, AE_{ij}) (suitably identified with column vectors in $\mathbb{R}^{2n^2}$), and let G be a $2n^2 \times 4n(n-1)$ matrix that has as its columns the vectors $(E_{ii} \pm E_{ij}, 0)$, $i \neq j$ and $(0, E_{ii} \pm E_{ij})$, $i \neq j$. Then $U \cap V$ is nonempty if and only if there exist vectors x and y , with $y > 0$, so that $Hx = Gy$. The components of y are positive because V is the interior of the cone generated by the columns of G . The problem of finding x and y , if they exist, is a linear program with $2n^2$ rows and $5n^2 - 4n$ columns.

Most linear program solvers require the inequality $y > 0$ to be of the form $\geq$. We will therefore require that the components of y each be at least 1. This does not affect the feasibility of the problem, because V is closed under multiplication by positive scalars. We also look for a solution that minimizes the sum of the components of y .

3. Infeasibility Proof. From Theorem 11.2 of Rockafellar [5] we conclude that if the subspace V and the non-empty open convex set U are disjoint, there must be a hyperplane L containing V so that one of the open half-spaces associated with L contains U . Clearly, U is in an open half-space associated with L if and only if its closure is in the corresponding closed half-space and an element of the closure of U is in the open half-space. When V is the span of the columns of a matrix H , as above, and U is the interior of the cone generated by the columns of a matrix G , the existence of such an L follows from the theorem of Stiemke [6].

The dot product of two elements (R, S) and (C, B) in $\mathbb{R}^{2n^2}$ is $(R, S) \cdot (C, B) = \#(R \circ C + S \circ B)$ where $\circ$ is the entrywise product and $\#$ is the sum of the entries function. Let L be the set of (C, B) in $\mathbb{R}^{2n^2}$ for which $(R, S) \cdot (C, B) = 0$. Then $V \subseteq L$ if and only if $(R, S) \cdot (E_{ij}, AE_{ij}) = 0$ for all pairs (i, j) . Note that AE_{ij} is the matrix with column i of A in its j^{th} column and zeros everywhere else. It follows

that $\#(R \circ E_{ij}) = r_{ij}$ and that $\#(S \circ AE_{ij})$ is the usual dot product of column j of S and column i of A . The inclusion $V \subseteq L$ is therefore equivalent to the equation $R + A^T S = 0$.

The requirement that the closure of U be in one of the closed half-spaces determined by L is equivalent to the set of $4n(n-1)$ inequalities $(R, S) \cdot (E_{ii} \pm E_{ij}, 0) \geq 0$, $i \neq j$ and $(R, S) \cdot (0, E_{ii} \pm E_{ij}) \geq 0$, $i \neq j$. The inequalities $(R, S) \cdot (E_{ii} \pm E_{ij}, 0) \geq 0$ for a given pair (i, j) with $i \neq j$ are clearly equivalent to the inequality $r_{ii} \geq |r_{ij}|$, and the inequalities $(R, S) \cdot (0, E_{ii} \pm E_{ij}) \geq 0$ for a given pair (i, j) with $i \neq j$ are clearly equivalent to the inequality $s_{ii} \geq |s_{ij}|$.

In order for the open set U , which is in a closed halfspace determined by L , to be in the corresponding open halfspace, we need one of the extreme rays of the closure of U to miss L . Thus, for some $i \neq j$, we need one of the following: $r_{ii} > r_{ij}$, $r_{ii} > -r_{ij}$, $s_{ii} > s_{ij}$, or $s_{ii} > -s_{ij}$. For a pair (R, S) that satisfies the inequalities $r_{ii} \geq |r_{ij}|$ and $s_{ii} \geq |s_{ij}|$ whenever $i \neq j$, the satisfaction of at least one such strict inequality is equivalent to $(R, S) \neq (0, 0)$.

It is not possible for $U \cap V$ to be nonempty if U is in one of the open half-spaces defined by L (which contains V). We therefore have the following theorem.

THEOREM 3.1. *Let A be an $n \times n$ matrix. Exactly one of the following holds:*

1. *There exist row diagonally dominant matrices B and C with positive diagonal entries and with $AC = B$, or*
2. *There exist square matrices R and S , not both 0, satisfying $R + A^T S = 0$ and, for all $1 \leq i, j \leq n$, $r_{ii} \geq |r_{ij}|$ and $s_{ii} \geq |s_{ij}|$.*

4. A non-hprdd matrix that is positive definite. We implemented the linear program of the previous section using the free software "MPL 4.2 for Windows." The 6×6 symmetric positive definite matrix of [1] that is not hidden M encouraged us to test symmetric positive definite matrices. (See [3] for a 3×3 positive definite matrix that is not hidden M .) It is also easy to generate random symmetric positive definite matrices by randomly generating their Cholesky factors. The smallest counterexample we found is the matrix

$$\begin{bmatrix} 4 & 4 & 7 & -9 \\ 4 & 16 & -7 & -2 \\ 7 & -7 & 30 & -24 \\ -9 & -2 & -24 & 25 \end{bmatrix},$$

for which the proof that it is not hprdd is provided by the matrices

$$R = \begin{bmatrix} 11 & 8 & 2 & -11 \\ 23 & 24 & -24 & -7 \\ 10 & -16 & 23 & -22 \\ -22 & -8 & -23 & 23 \end{bmatrix} \text{ and } S = \begin{bmatrix} 24 & -24 & -9 & 23 \\ -8 & 8 & 8 & -5 \\ -3 & 8 & 8 & -1 \\ 6 & 0 & 6 & 6 \end{bmatrix}.$$

We should point out that this approach does not resolve the question of whether or not there exist 3×3 P -matrices satisfying the second alternative of the Theorem. To produce a P matrix satisfying the second alternative where its transpose does not is an interesting challenge.

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