BOUNDS FOR THE COMPLETELY POSITIVE RANK OF A SYMMETRIC MATRIX OVER A TROPICAL SEMIRING*

DAVID DOLŽAN[†] AND POLONA OBLAK[‡]

Abstract. In this paper, an upper bound for the CP-rank of a matrix over a tropical semiring is obtained, according to the vertex clique cover of the graph prescribed by the positions of zero entries in the matrix. The graphs that beget the matrices with the lowest possible CP-ranks are studied, and it is proved that any such graph must have its diameter equal to 2.

Key words. Tropical semiring, Symmetric matrix, Rank.

AMS subject classifications. 15A23, 15B48, 16Y60.

1. Introduction. In this paper, we study the completely positive rank of a matrix over the tropical semiring \mathbb{T} , which is the semiring $(\mathbb{R} \cup \{\infty\}, \oplus, \odot)$, with operations defined by $a \oplus b = \min\{a, b\}$ and $a \odot b = a + b$.

For a semiring S, we say that a symmetric $n \times n$ matrix A over S is completely positive, if there exists an $n \times r$ matrix B over S such that

$$A = BB^T$$
.

The minimal possible r in such factorization, is the *CP*-rank of A and it is denoted by CPrk(A). Equivalently, a matrix A has CPrk(A) = r if and only if r is the smallest number, such that there exist vectors $b_1, b_2, \ldots, b_r \in \mathbb{T}^n$ with

$$A = \sum_{i=1}^{r} b_i b_i^T.$$

If matrix A is not completely positive, we denote $\operatorname{CPrk}(A) = \infty$. Note that in [6], the authors refer to CP-rank as the symmetric Barvinok rank of a matrix.

Note that over semirings, all definitions of the rank of a matrix do not coincide as in the case of matrices over real numbers with standard operations (see e.g. [1, 9]). Thus, the CP-rank (which is a special case of a factor rank) is just one of many possible semiring matrix ranks.

For a completely positive $n \times n$ matrix A over the field \mathbb{R} , Drew, Johnson and Loewy [7] conjectured that $\operatorname{CPrk}(A) \leq \lfloor \frac{n^2}{4} \rfloor$ if $n \geq 4$. Twenty years later, the conjectured upper bound was proved wrong and corrected to $\frac{n^2}{2}$ for all $n \geq 7$ [4, 5]. However, it is still not known what is the tight upper bound and it transpires that the problem of determining the CP-rank of any given matrix is a difficult problem [2, 3].

^{*}Received by the editors on August 9, 2017. Accepted for publication on March 13, 2018. Handling Editor: Sergey Sergeev. Corresponding Author: Polona Oblak. The authors acknowledge the financial support from the Slovenian Research Agency (research core funding no. P1-0222).

[†]Department of Mathematics, Faculty of Mathematics and Physics, University of Ljubljana, Jadranska 21, SI-1000 Ljubljana, Slovenia (david.dolzan@fmf.uni-lj.si).

[‡]Faculty of Computer and Information Science, University of Ljubljana, Večna pot 113, SI-1000 Ljubljana, Slovenia (polona.oblak@fri.uni-lj.si).

Let $M_n(S)$ denote the semiring of all $n \times n$ matrices over the semiring S. Over the tropical semiring \mathbb{T} , Cartwright and Chan [6] proved that $\max\left\{n, \left\lfloor \frac{n^2}{4} \right\rfloor\right\}$ is the tight upper bound for the CP-rank of a completely positive matrix $A \in M_n(\mathbb{T})$. Over the Boolean semiring and the max-min semiring, the same inequality was proved by Mohindru [11] and Shitov [14].

In [13], Shaked-Monderer introduced CPrk (G) to be the maximum CP-rank of all real matrices with the pattern prescribed by the graph G. She proved that the CPrk (G) is equal to the to the edge clique cover number of G, if and only if G is not a tree and does not contain a triangle.

We follow [6] to define $\operatorname{CPrk}(G)$ over the tropical semiring to be the maximum of $\operatorname{CP-ranks}$ of all completely positive matrices $A = (a_{ij}) \in M_n(\mathbb{T})$ such that, for $i \neq j$, $a_{ij} = 0$ if and only if $\{i, j\} \in E(G)$. (Note that throughout the paper, zero is a real number and not the tropical additive identity, which is ∞ .) Observe that in G, edges correspond to all entries equal to a specific element 0 distinct from the additive identity in \mathbb{T} . This graph is a subgraph of the weighted graph corresponding to a semiring matrix (see for example [8]), which is also called the precedence graph.

In this paper, we find an upper bound for the CP-rank of a matrix with regards to the vertex clique cover of the graph prescribed by the positions of zero entries in the matrix. This bound can be much lower than the bound max $\left\{n, \left\lfloor \frac{n^2}{4} \right\rfloor\right\}$ from [6, Theorem 4], see Theorem 3.4 and Remark 3.5. We then proceed to apply these results to 0/1 matrices, since it was established in [6] that CP-rank of 0/1 matrices is equal to the edge clique cover number of the corresponding graph. We examine the connection between the ranks of 0/1 matrices and arbitrary matrices with the same positions of zero entries. In the last section, we then study the graphs that beget the matrices with the lowest possible CP-ranks. We prove that any such graph must have its diameter equal to 2, and provide examples that in case of diameter 2 the rank does not seem to be well behaved.

2. Preliminary results. In this section, we give the basic definitions and some preliminary results. First, we provide the characterization of completely positive matrices over the tropical semiring. The subset of $M_n(\mathbb{T})$ of all completely positive matrices will be denoted by $\operatorname{CP}_n(\mathbb{T})$.

The following lemma is obvious and characterizes matrices of CP-rank equal to 1.

LEMMA 2.1. A symmetric matrix $A = (a_{ij}) \in M_n(\mathbb{T})$ has $\operatorname{CPrk}(A) = 1$ if and only if $a_{ij_1} \odot a_{kj_2} = a_{kj_1} \odot a_{ij_2}$ for all $i, j_1, j_2, k = 1, 2, \ldots, n$. (This means that the difference between any two rows of A with finite entries is a vector with all of its entries equal.)

The following lemma characterizes completely positive matrices over the tropical semiring.

LEMMA 2.2. [6, Proposition 2 and Theorem 4] A symmetric matrix $A = (a_{ij}) \in M_n(\mathbb{T})$ is completely positive if and only if $2a_{ij} \ge a_{ii} + a_{jj}$ for all i, j = 1, 2, ..., n.

This lemma implies that if $a_{ii} = \infty$ for $A = (a_{ij}) \in CP_n(\mathbb{T})$ and some *i*, then $a_{ij} = \infty$ for all *j*. Also, if all the diagonal elements of a completely positive matrix *A* are equal to 0, then all off-diagonal entries are nonnegative. This fact makes it convenient to study such matrices, and also gives sense to studying matrices defined by the positions of the zero entries. The next paragraph describes the procedure to transform the completely positive matrix with diagonal entries equal to 0, while preserving the CP-rank.

Choose $A = (a_{ij}) \in M_n(\mathbb{T})$. Let $A[i] \in M_{n-1}(\mathbb{T})$ be the matrix obtained from A by deleting its *i*-th row and *i*-th column and let $b[i] \in \mathbb{T}^{n-1}$ be the vector obtained from vector $b \in \mathbb{T}^n$ by deleting its *i*-th entry. If

David Dolžan and Polona Oblak

matrix A has k diagonal entries equal to ∞ , let $C(A) \in M_{n-k}(\mathbb{T})$ be the matrix obtained from A by

- deleting *i*-th row and *i*-th column if $a_{ii} = \infty$ for every i = 1, 2, ..., n, and
- subtracting $\frac{1}{2}a_{ii}$ from each entry in the *i*-th row and *i*-th column of A, if $a_{ii} \neq \infty$ for every $i = 1, 2, \ldots, n$. (Note that subtracting a real number from ∞ yields ∞ and that we subtract $\frac{1}{2}a_{ii}$ twice from a_{ii} .)

The next lemma assures us that the rank of a matrix does not change with the above transformation. LEMMA 2.3. If $A \in M_n(\mathbb{T})$ is completely positive, then

$$\operatorname{CPrk}(A) = \operatorname{CPrk}(C(A)).$$

Proof. Let $A = (a_{ij}) = \bigoplus_{j=1}^r b_j \odot b_j^T \in CP_n(\mathbb{T})$ and suppose first that $a_{ii} = \infty$ for some $i, 1 \leq i \leq n$. Observe that $A[i] = \bigoplus_{j=1}^r b_j[i] \odot b_j[i]^T \in CP_{n-1}(\mathbb{T})$, which implies that $CPrk(A) \geq CPrk(A[i])$. Similarly, we can observe that $CPrk(A[i]) \geq CPrk(A)$, by inserting a component equal to ∞ to all b_j at the *i*-th component, since a completely positive matrix A with $a_{ii} = \infty$, by Lemma 2.2 must have all entries in the *i*-th row and *i*-th column equal to ∞ .

Now, suppose $A = \bigoplus_{j=1}^{r} b_j \odot b_j^T \in CP_n(\mathbb{T})$ and $a_{ii} \neq \infty$ for i = 1, 2..., n. Choose $\alpha \in \mathbb{R}, k \in \{1, 2, ..., n\}$, and let $B \in CP_n(\mathbb{T})$ and $c_j \in \mathbb{T}^n$ be defined as

$$B_{ij} = \begin{cases} a_{ij} + 2\alpha, & \text{if } i = j = k, \\ a_{ij} + \alpha, & \text{if either } i = k \text{ or } j = k, \\ a_{ij}, & \text{otherwise,} \end{cases} \text{ and } (c_j)_i = \begin{cases} (b_j)_i + \alpha, & \text{if } i = k, \\ (b_j)_i, & \text{otherwise.} \end{cases}$$

Observe that $B = \bigoplus_{j=1}^{r} c_j \odot c_j^T$, and thus, $\operatorname{CPrk}(B) \leq \operatorname{CPrk}(A)$. By replacing α by $-\alpha$, we obtain $\operatorname{CPrk}(A) \leq \operatorname{CPrk}(B)$. By consecutively applying the above procedure with $\alpha = -\frac{1}{2}a_{kk}$ for all $k = 1, 2, \ldots, n$, we conclude that $\operatorname{CPrk}(A) = \operatorname{CPrk}(C(A))$.

The next example shows that in general, the positions of nonzero entries in a matrix do not determine the CP-rank. We shall see later that this inconvenience can be circumnavigated by replacing A with C(A)as described above, which is a transformation that preserves the CP-rank by Lemma 2.3.

EXAMPLE 2.4. Let

$$A = \begin{bmatrix} 0 & 1 & 2 \\ 1 & 2 & 3 \\ 2 & 3 & 4 \end{bmatrix} = \begin{bmatrix} 0 \\ 1 \\ 2 \end{bmatrix} \odot \begin{bmatrix} 0 & 1 & 2 \end{bmatrix} \in \operatorname{CP}_3(\mathbb{T})$$

By transformation described on page 154, we obtain

$$C(A) = \begin{bmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} \odot \begin{bmatrix} 0 & 0 & 0 \end{bmatrix}$$

and have $\operatorname{CPrk}(A) = \operatorname{CPrk}(C(A)) = 1$.

Note that by changing the nonzero entries of matrix A, we obtain a matrix with different CP-rank. For example, if

$$B = \begin{bmatrix} 0 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{bmatrix},$$



then

$$C(B) = \begin{bmatrix} 0 & \frac{1}{2} & \frac{1}{2} \\ \frac{1}{2} & 0 & 0 \\ \frac{1}{2} & 0 & 0 \end{bmatrix} = \begin{bmatrix} 0 \\ \frac{1}{2} \\ \frac{1}{2} \end{bmatrix} \odot \begin{bmatrix} 0 & \frac{1}{2} & \frac{1}{2} \end{bmatrix} \oplus \begin{bmatrix} \infty \\ 0 \\ 0 \end{bmatrix} \odot \begin{bmatrix} \infty & 0 & 0 \end{bmatrix}$$

Lemma 2.1 implies that $\operatorname{CPrk}(B) \neq 1$. Note that $\operatorname{CPrk}(C(B)) \leq 2$ and by Lemma 2.3 it follows that $\operatorname{CPrk}(B) = \operatorname{CPrk}(C(B)) = 2$.

3. Bounding the CP-rank by the graph structure. In this section, we find bounds for CP-ranks of matrices with the aid of a graph structure that is prescribed to a given matrix. Namely, we define a graph that corresponds to a matrix (depending on whether different elements of the matrix are equal to zero). We find bounds for the CP-rank of all matrices with a given graph structure. Note that using Lemma 2.3, we always work under the assumption that $A \in CP_n(\mathbb{T})$ has a zero diagonal and nonnegative offdiagonal entries.

Given a symmetric matrix $A = (a_{ij}) \in M_n(\mathbb{T})$, we define G(A) = (V, E) to be a simple graph with $V = \{1, 2, ..., n\}$, and for $i \neq j$ we have $\{i, j\} \in E$ if and only if $a_{ij} = 0$. Recall that CPrk (G) is the maximum of CP-ranks of all symmetric matrices $A = (a_{ij}) \in M_n(\mathbb{T})$ such that, for $i \neq j$, $a_{ij} = 0$ if and only if $\{i, j\} \in E(G)$.

As usual, in a given graph, the path $a = x_0 \sim x_1 \sim \cdots \sim x_{n-1} \sim x_n = b$ connecting vertices a and b has length n and the length of the shortest path connecting vertices a and b is called the distance between a and b and denoted by d(a, b). We let $d(a, b) = \infty$ if there is no path connecting a and b, and we let d(a, a) = 0. The diameter of a graph is a maximal distance between any two of its vertices. An empty graph is a graph consisting of isolated nodes with no edges. A complete graph on n vertices will be denoted by K_n and a path with n vertices will be denoted by P_n . The edge clique cover number cc(G) of a graph G is the minimal cardinality of the collections of complete subgraphs such that every edge of G is in one element of the collection.

The following two lemmas give us some bounds for the CP-rank of graphs and their subgraphs.

LEMMA 3.1. If H is an induced subgraph of the graph G, then

$$\operatorname{CPrk}(H) \leq \operatorname{CPrk}(G)$$
.

Proof. Let H be an induced subgraph of G and suppose without loss of generality that $V(H) = \{1, 2, ..., m\}$ and $V(G) = \{1, 2, ..., n\}$, $m \leq n$. Choose any $A \in M_n(\mathbb{T})$ with G(A) = G, and let B be its $m \times m$ leading principal submatrix. It is clear that G(B) = H. If $A = \bigoplus_{i=1}^k a_i \odot a_i^T$, then $B = \bigoplus_{i=1}^k b_i \odot b_i^T$, where b_i is a vector obtained from a_i by deleting its last n - m components. Hence, $\operatorname{CPrk}(B) \leq \operatorname{CPrk}(A)$ and so $\operatorname{CPrk}(H) \leq \operatorname{CPrk}(G)$.

Recall that the join $G \vee H$ of graphs G and H, is the graph union $G \cup H$ together with all the possible edges joining the vertices in G to the vertices in H. We show in the next lemma that joining a graph with a single vertex does not change the CP-rank.

LEMMA 3.2. For any graph G and w a vertex not in G, we have

$$\operatorname{CPrk}(G \lor w) = \operatorname{CPrk}(G).$$

Proof. Take any $A \in M_{n+1}(\mathbb{T})$ with $G(A) = G \vee w$. Hence, A is a direct sum of matrix $B \in M_n(\mathbb{T})$, G(B) = G, with the size one zero matrix. There exist $b_i \in \mathbb{T}^n$, $i = 1, 2, \ldots, k, k \leq \operatorname{CPrk}(G)$, such that

I L
AS

David Dolžan and Polona Oblak

 $B = \bigoplus_{i=1}^{k} b_i \odot b_i^T. \text{ Define } a_i = \begin{bmatrix} b_i^T & 0 \end{bmatrix}^T \in \mathbb{T}^{n+1} \text{ and observe that } A = \bigoplus_{i=1}^{k} a_i \odot a_i^T \text{ and hence, } \operatorname{CPrk}(A) \leq k.$ This implies that $\operatorname{CPrk}(G \lor w) \leq \operatorname{CPrk}(G).$ By Lemma 3.1, it follows that $\operatorname{CPrk}(G \lor w) = \operatorname{CPrk}(G).$

Now, we define the vertex clique cover γ of a graph G as a collection of r complete subgraphs such that every vertex of G is in some element of the collection. One can always assume that the vertices of G are labeled so that

$$\gamma = (K_{q_1}, K_{q_2}, \dots, K_{q_k}, \underbrace{K_1, \dots, K_1}_l) = (K_{q_1}, K_{q_2}, \dots, K_{q_k}, lK_1),$$

where $q_1 \ge q_2 \ge \cdots \ge q_k \ge 2$. Define the vertex clique cover number of γ as

$$\theta(\gamma) = k + \sum_{i=1}^{k} (i-1)q_i + kl + \left\lfloor \frac{l^2}{4} \right\rfloor.$$

It is worth noting that a vertex clique cover number is the same as a chromatic number of the complement of the graph. In Theorem 3.4, we will prove that CP-rank of a matrix A is bounded by $\theta(\gamma)$ for any vertex clique cover γ of G = G(A).

EXAMPLE 3.3. Note that a vertex clique cover is not unique. Let G be a paw graph:



Its vertex clique covers are

$$\gamma_1 = (K_3, K_1)$$
 and $\gamma_2 = (2K_2)$,

so $\theta(\gamma_1) = 2$ and $\theta(\gamma_2) = 4$.

The next theorem specifies an upper bound for the CP-rank of a matrix according to the vertex clique cover number of a graph corresponding to the matrix.

THEOREM 3.4. Choose $A \in CP_n(\mathbb{T})$. If G(A) is a nonempty graph or $n \geq 5$, then for every vertex clique cover γ of G(A), we have

$$\operatorname{CPrk}(A) \leq \theta(\gamma).$$

Otherwise, if G(A) is an empty graph with $n \leq 4$, then

$$\operatorname{CPrk}(A) = n.$$

Proof. Suppose first that G = G(A) is nonempty graph. We will construct $n \times n$ matrices A_1, A_2, A_3 and $A_4, A = A_1 \oplus A_2 \oplus A_3 \oplus A_4$, which will correspond to subgraphs of G, and their CP-ranks will be bounded by $k, \sum_{i=1}^{k} (i-1)q_i, kl$ and $\lfloor \frac{l^2}{4} \rfloor$, respectively.

1. If k = 0, then A_1 is a zero matrix. Suppose $k \ge 1$. For i = 1, 2, ..., k denote the components of $x^{(i)} \in \mathbb{T}^n$ by

$$x_{j}^{(i)} = \begin{cases} 0, & \text{if } q_{1} + \dots + q_{i-1} + 1 \leq j \leq q_{1} + \dots + q_{i-1} + q_{i} \\ \infty, & \text{otherwise,} \end{cases}$$



for all $j = 1, 2, \ldots, n$. Define

$$A_1 = \bigoplus_{i=1}^k x^{(i)} \odot \left(x^{(i)}\right)^T$$

is a sum of k matrices of CP-rank one. Note that A_1 coincides with A at all elements that correspond to the edges of cliques K_{q_1} to K_{q_k} of G.

2. If $k \leq 1$, then A_2 is a zero matrix. Suppose that $k \geq 2$. For i = 1, 2, ..., k - 1, j = i + 1, i + 2, ..., kand $s = 1, 2, ..., q_j$ denote the components of $y^{(i,j,s)} \in \mathbb{T}^n$ by

$$y_t^{(i,j,s)} = \begin{cases} 0, & \text{if } t = q_1 + \dots + q_{j-1} + s, \\ a_{t,q_1 + \dots + q_{j-1} + s}, & \text{if } q_1 + \dots + q_{i-1} + 1 \le t \le q_1 + \dots + q_i, \\ \infty, & \text{otherwise.} \end{cases}$$

Let us define

$$A_2 = \bigoplus_{i=1j}^{k-1} \bigoplus_{i=i+1}^k \bigoplus_{s=1}^{q_j} y^{(i,j,s)} \odot \left(y^{(i,j,s)} \right)^T.$$

Note that A_2 is a sum of $\sum_{j=1}^{\kappa} (j-1)q_j$ matrices of CP-rank one, that coincides with the matrix A at

all elements that correspond to the edges between any of the cliques K_{q_1} to K_{q_k} of G.

3. If k = 0 or l = 0, then A_3 is a zero matrix. Suppose that $k, l \ge 1$. For i = 1, ..., k and j = 1, ..., l denote the components of $z^{(i,j)} \in \mathbb{T}^n$ by

$$z_t^{(i,j)} = \begin{cases} 0, & \text{if } t = q_1 + \dots + q_k + j, \\ a_{t,q_1 + \dots + q_k + j}, & \text{if } q_1 + \dots + q_{i-1} + 1 \le t \le q_1 + \dots + q_i, \\ \infty, & \text{otherwise.} \end{cases}$$

Let

$$A_3 = \bigoplus_{i=1}^k \bigoplus_{j=1}^l z^{(i,j)} \odot \left(z^{(i,j)} \right)^T$$

be a matrix defined as a sum of kl matrices of CP-rank one. Note that A_3 coincides with the matrix A at all elements that correspond to the edges between any of the clique K_1 and any of the cliques K_{q_1} to K_{q_k} of G.

4. If $l \leq 1$, then A_4 is a zero matrix and for $l \geq 4$ let the matrix A_4 be defined by

$$(A_4)_{ij} = \begin{cases} \infty, & \text{if } i \le q_1 + \dots + q_k \text{ or } j \le q_1 + \dots + q_k, \\ a_{ij}, & \text{otherwise.} \end{cases}$$

Note that A_4 coincides with the matrix A at all elements that correspond to the edges between any of the cliques K_1 of G.

If $l \ge 4$, then note that A_4 can be written as a sum of at most $\left\lfloor \frac{l^2}{4} \right\rfloor$ CP-rank one matrices by [6, Theorem 4].

In the case $2 \leq l \leq 3$, observe that $n \geq 5$ implies that k > 0. This further implies that $A_3 \neq 0$, and thus, $(A_3)_{ii} = 0$ for $i \geq q_1 + \cdots + q_k + 1$, by the construction of A_3 above. For l = 2, matrix A_4 is of CP-rank $\left\lfloor \frac{2^2}{4} \right\rfloor = 1$, since $A_4 = [\infty, \dots, \infty, 0, a_{n-1,n}]^T \odot [\infty, \dots, \infty, 0, a_{n-1,n}]$. For l = 3, assume without loss of generality that $a_{n-1,n} = \max\{a_{n-2,n-1}, a_{n-2,n}, a_{n-1,n}\}$. In this case, we have $A_4 = a \odot a^T \oplus b \odot b^T$, where $a = \begin{bmatrix} \infty \cdots \infty & 0 & a_{n-2,n-1} & \infty \end{bmatrix}^T \in \mathbb{T}^n$ and $b = \begin{bmatrix} \infty \cdots \infty & a_{n-2,n} & a_{n-1,n} & 0 \end{bmatrix}^T \in \mathbb{T}^n$. It follows that $\operatorname{CPrk}(A_4) = 2 = \left\lfloor \frac{3^2}{4} \right\rfloor$.

David Dolžan and Polona Oblak

Observe that

$$A = A_1 \oplus A_2 \oplus A_3 \oplus A_4,$$

and therefore, the inequality in the statement follows.

If G is an empty graph, then k = 0. In addition, if $n = l \ge 5$, we construct A_4 as above, and then $A = A_4$ is a sum of at most $\lfloor \frac{n^2}{4} \rfloor$ matrices of CP rank one. If $n \le 4$, then observe that $\lfloor \frac{n^2}{4} \rfloor \le n$, so by [6, Theorem 4] A can be written as a sum of at most n matrices of CP rank one. However, since G is an empty graph, each summand with CP rank one can have at most one zero element. Since $A = A_4$ has zeroes on the diagonal, this implies that there must be exactly n summands with CP rank one.

REMARK 3.5. Note that $\theta(\gamma)$ is a much smaller number than $\left\lfloor \frac{n^2}{4} \right\rfloor$ whenever $k \ge 1$, so there are infinite families of graphs and consequently infinite families of matrices for which we have found a much lower bound for their CP rank. For example, when k = 1 (and similarly, one can reason for all other $k \ge 1$), $\theta(\gamma) = 1 + l + \left\lfloor \frac{l^2}{4} \right\rfloor$, which (since q_1 can be arbitrarly large) can actually be arbitrarily smaller than $\left\lfloor \frac{(q_1+l)^2}{4} \right\rfloor$.

EXAMPLE 3.6. Theorem 3.4 implies that any matrix

$$A = \begin{bmatrix} 0 & 0 & 0 & a \\ 0 & 0 & 0 & b \\ 0 & 0 & 0 & 0 \\ a & b & 0 & 0 \end{bmatrix} \in \operatorname{CP}_4(\mathbb{T}),$$

where a, b > 0, which corresponds to the paw graph from Example 3.6, has $\operatorname{CPrk}(A) \leq \theta(\gamma_1) < \theta(\gamma_2)$. Note that by Lemma 2.1 it follows that $\operatorname{CPrk}(A) = 2$.

The next example shows that CP-rank of a matrix A with an empty graph can be strictly greater than n, when n > 4.

EXAMPLE 3.7. Let

$$A = \begin{bmatrix} 0 & 1 & 1 & 3 & 3 \\ 1 & 0 & 3 & 1 & 1 \\ 1 & 3 & 0 & 1 & 1 \\ 3 & 1 & 1 & 0 & 3 \\ 3 & 1 & 1 & 3 & 0 \end{bmatrix} \in \operatorname{CP}_{5}(\mathbb{T}),$$

and let us prove that $\operatorname{CPrk}(A) = 6$.

Suppose there exist vectors $b_1, b_2, \ldots, b_5 \in \mathbb{T}^5$ such that

$$A = \bigoplus_{i=1}^{5} b_i \odot b_i^T$$

Since all diagonal entries of $A = (a_{ij})$ are equal to zero and all offdiagonal entries are nonzero, it follows that each $b_i = [b_{i1}, b_{i2}, \ldots, b_{i5}]^T$ has nonnegative entries with exactly one zero entry. Without any loss of generality, we asume that $b_{ii} = 0, i = 1, 2, \ldots, 5$.

Let us define $\mathcal{E} = \{(1,2), (1,3), (2,4), (2,5), (3,4), (3,5)\}$ the set of indices such that, for k < l, we have $a_{kl} = 1$ if and only if $(k,l) \in \mathcal{E}$. Note that $b_{kl} = b_{kk} + b_{kl} \ge a_{kl}$ for all $1 \le k < l \le 5$, which gives us $b_{kl} \ge 1$ for $(k,l) \in \mathcal{E}$ and $b_{kl} \ge 3$ for $(k,l) \notin \mathcal{E}$.

158

Moreover, for any pair $(k,l) \notin \mathcal{E}$, $1 \leq k < l \leq 5$, and any *i*, we have $b_{ki} + b_{il} \geq a_{kl} = 3$. This gives us

	$b_{21} + b_{24} \ge 3,$	$b_{31} + b_{34} \ge 3,$	$b_{21} + b_{25} \ge 3,$
(3.1)	$b_{31} + b_{35} \ge 3,$	$b_{12} + b_{13} \ge 3,$	$b_{42} + b_{43} \ge 3,$
	$b_{52} + b_{53} \ge 3,$	$b_{24} + b_{25} \ge 3,$	$b_{34} + b_{35} \ge 3.$

Note that $b_{ik} + b_{il} \ge 2$ for all *i* distinct from *k* and *l*, and thus,

(3.2)
$$\min\{b_{kl}, b_{lk}\} = 1$$

for any $(k, l) \in \mathcal{E}$.

Choose (k, l) = (1, 2) and by (3.2) we have $b_{21} = 1$ or $b_{12} = 1$. In the case $b_{21} = 1$, we apply (3.1) and (3.2) for several times, to observe that $b_{24} \ge 2$, $b_{25} \ge 2$, $b_{42} = 1$, $b_{43} \ge 2$, $b_{34} = 1$, $b_{35} \ge 2$, $b_{53} = 1$, $b_{52} \ge 2$ and so $b_{25} = 1$, a contradiction. Similar arguments give us a contradiction also in the case $b_{12} = 1$. Hence, we proved that CPrk $(A) \ge 6$ and by Theorem 3.4, it follows that CPrk (A) = 6.

In the rest of this section, we apply the above results to the study of the CP-rank of 0/1 matrices over \mathbb{T} . Note again that 0 and 1 here represent real numbers. Equivalently, one could also study $0/\infty$ matrices, where 0 and ∞ represent the tropical identity and tropical zero.

It can be seen that CP-rank of a 0/1 matrix A is equal to the edge clique cover number of G(A), denoted by cc (G(A)) [6, Proposition 3]. Note that it was proved that the edge clique cover number of a graph is equal to the intersection number of the graph [10]. Since determining the intersection number is an NP-complete problem [12], it seems useful to obtain some easily calculable bounds for the CP-rank of a 0/1 matrix and the following two propositions offer some results in this direction, by using the same approach as in the proof of Theorem 3.4.

PROPOSITION 3.8. If $A \in CP_n(\mathbb{T})$ is a 0/1 matrix such that G(A) is an empty graph, then

$$\operatorname{CPrk}(A) = n.$$

Proof. Let us define $v^{(i)} \in \mathbb{T}^n$, $i = 1, 2, \ldots, n$, by

$$v_j^{(i)} = \begin{cases} 0, & \text{if } i = j, \\ 1, & \text{if } i \neq j. \end{cases}$$

It is easy to verify that

$$A = \bigoplus_{i=1}^{n} v^{(i)} \odot \left(v^{(i)} \right)^{T}$$

and so $\operatorname{CPrk}(A) \leq n$. Suppose now $A = \bigoplus_{i=1}^{n-1} u^{(i)} \odot (u^{(i)})^T$. Since A has n diagonal entries equal to 0, there exists j such that $u_t^{(j)} = u_s^{(j)} = 0$ for some $1 \leq s, t \leq n$. It follows that $(u^{(i)} \odot (u^{(i)})^T)_{ts} = 0$, and thus, $a_{ts} \neq 0$, a contradiction. Therefore, $\operatorname{CPrk}(A) = n$.

Note that the above proposition is not valid for matrices which are not 0/1, as Example 3.7 shows.

For any given matrix $A \in M_n(\mathbb{T})$, we define its support, $\text{Supp}(A) \in M_n(\mathbb{T})$, by

$$\operatorname{Supp}(A)_{ij} = \begin{cases} 0, & \text{if } a_{ij} = 0, \\ 1, & \text{if } a_{ij} \neq 0. \end{cases}$$

David Dolžan and Polona Oblak

In Example 2.4, we showed that the CP-rank of A and Supp(A) do not necessarily coincide.

LEMMA 3.9. If G is a graph with $\operatorname{CPrk}(G) = \operatorname{cc}(G)$, then for every $A = (a_{ij}) \in \operatorname{CP}_n(\mathbb{T})$ with G(A) = G choose edge clique cover $Q_1, Q_2, \ldots, Q_{\operatorname{cc}(G)}$. Then

(3.3)
$$A = \bigoplus_{i=1}^{\operatorname{cc}(G)} b_i \odot b_i^T$$

and the following two statements hold:

- (a) We have a bijective correspondence between the cliques $Q_1, Q_2, \ldots, Q_{cc(G)}$ and the summands b_i of the sum, where the vertices of the clique *i* correspond to the zero entries of b_i .
- (b) If a_{uv} is the minimal nonzero entry in A, then for every i = 1, 2, ..., cc(G) and j = 1, 2, ..., n, we have $(b_i)_j = 0$ or $(b_i)_j \ge a_{u,v}$.

Proof. Since G(A) = G and $\operatorname{CPrk}(G) = \operatorname{cc}(G)$, we know that $A = \bigoplus_{i=1}^{\operatorname{cc}(G)} b_i \odot b_i^T$ for some vectors $b_i \in \mathbb{T}^n$. For every clique Q from the clique cover $Q_1, Q_2, \ldots, Q_{\operatorname{cc}(G)}$, we have $a_{jk} = 0$ for all $j, k \in Q$. This implies that there exists i such that $(b_i)_j = (b_i)_k = 0$ for all $j, k \in Q$. The fact that the number of summands of rank one matrices is exactly equal to $\operatorname{cc}(G)$, implies that for every clique Q_i in G, there exists some vector b_i with components equaling zero at least at all positions corresponding to the vertices of clique Q_i . By Lemma 2.3 and the definition of operations in $M_n(\mathbb{T})$, we know that all positions that correspond to vertices outside clique Q_i , have to be nonzero. This yields the desired bijective correspondence.

Now, suppose a_{uv} is the minimal nonzero entry in A and choose i such that $(b_i)_j > 0$. By the above, b_i corresponds to a clique Q_i in G, so there exist indices k_1, k_2, \ldots, k_r such that $(b_i)_{k_t} = 0$ for all $t = 1, 2, \ldots, r$ and $j \notin Q_i$. Then $a_{jk_t} \leq (b_i \odot b_i^T)_{j,k_t} = (b_i)_j$ for all $t = 1, 2, \ldots, r$. Since vertices corresponding to j and k_t do not belong to the same clique, there exists at least one t such that $a_{jk_t} \neq 0$, and therefore $a_{uv} \leq a_{jk_t} \leq (b_i)_j$.

By [6, Proposition 3], we have that the CP-rank of Supp(A), which is a 0/1 matrix, is equal to the edge clique cover number of G(A). Therefore it follows that in order to find lower bounds for the CP-rank of any matrix, it suffices to study the CP-rank of its corresponding support as the following shows.

COROLLARY 3.10. For any matrix $A \in CP_n(\mathbb{T})$, we have

 $\operatorname{CPrk}(G(A)) \ge \operatorname{CPrk}(A) \ge \operatorname{CPrk}(\operatorname{Supp}(A)) = \operatorname{cc}(G(A)).$

Note that Example 2.4 shows that the inequality in Corollary 3.10 is not necessarily true if we omit the condition A = C(A).

4. Graphs with CP-rank equal to the clique cover number. In Lemma 3.9, we proved that the lower bound for CP-rank of a graph is its clique cover number. Therefore, we now proceed by studying the graphs that define matrices with the CP-ranks that are as close as possible to the bound from Corollary 3.10.

The following theorem shows that if we aspire to characterize graphs with the lowest possible CP-ranks, we can limit ourselves to graphs which are very well connected, i.e. their diameters are at most 2. However, the situation in the case diam $(G) \leq 2$ appears to be quite complex. We provide examples of acyclic and cyclic graphs with diameter 2 where either CPrk (G) = cc(G) or CPrk (G) > cc(G).

THEOREM 4.1. If G is a connected graph with $\operatorname{CPrk}(G) = \operatorname{cc}(G)$, then $\operatorname{diam}(G) \leq 2$.

160



Proof. Suppose $\operatorname{CPrk}(G) = \operatorname{cc}(G)$ and $\operatorname{diam}(G) \geq 3$. Thus, there exist vertices $u, v \in V(G)$ with $d(u, v) \geq 3$.

Define $A = (a_{ij}) \in M_n(\mathbb{T})$ by

$$a_{ij} = \begin{cases} 0, & \text{if } \{i, j\} \in E(G) \text{ or } i = j, \\ 1, & \text{if } \{i, j\} = \{u, v\}, \\ 2, & \text{if } \{i, j\} \notin E(G) \text{ and } i \neq j. \end{cases}$$

Observe that G(A) = G and $A \in CP_n(\mathbb{T})$. Let A be of the form (3.3). Since $a_{u,v} = 1$ is the minimal nonzero entry of A and $(b_i \odot b_i^T)_{u,v} = (b_i)_u + (b_i)_v = 1$ for some i, then by Lemma 3.9 (b), $(b_i)_u = 0$ and $(b_i)_v = 1$ or $(b_i)_v = 0$ and $(b_i)_u = 1$. Suppose without loss of generality that $(b_i)_u = 0$ and $(b_i)_v = 1$. By Lemma 3.9 (a), $(b_i)_l = 0$ for some $l \neq u$, and thus, $a_{v,l} \leq (b_i \odot b_i^T)_{v,l} = (b_i)_v + (b_i)_l = 1$. Hence, by definition of A, $\{v, l\} \in E(G)$, which contradicts $d(u, v) \geq 3$.

EXAMPLE 4.2. If $G = P_3$ is a path on 3 vertices, then all matrices $A \in CP_3(\mathbb{T})$ with $G(A) = P_3$ have (up to a permutational conjugation) the form

$$A = \begin{bmatrix} 0 & a & 0 \\ a & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix} = \begin{bmatrix} a \\ 0 \\ 0 \end{bmatrix} \odot \begin{bmatrix} a & 0 & 0 \end{bmatrix} \oplus \begin{bmatrix} 0 \\ \infty \\ 0 \end{bmatrix} \odot \begin{bmatrix} 0 & \infty & 0 \end{bmatrix}$$

for some $0 \neq a \in \mathbb{T}$. By Lemma 2.1, $\operatorname{CPrk}(A) \neq 1$, so it follows that $\operatorname{CPrk}(A) = 2$, and thus, $\operatorname{CPrk}(P_3) = \operatorname{cc}(P_3) = 2$.

EXAMPLE 4.3. If G is a paw graph (see Example 3.6), then cc(G) = 2. Since every matrix $B \in CP_4(\mathbb{T})$, G(B) = G, has (up to permutational conjugation) the form

$$B = \begin{bmatrix} 0 & a & b & 0 \\ a & 0 & 0 & 0 \\ b & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} = \begin{bmatrix} \infty \\ 0 \\ 0 \\ 0 \\ 0 \end{bmatrix} \odot \begin{bmatrix} \infty & 0 & 0 & 0 \end{bmatrix} \oplus \begin{bmatrix} 0 \\ a \\ b \\ 0 \end{bmatrix} \odot \begin{bmatrix} 0 & a & b & 0 \end{bmatrix}$$

for some $0 \neq a, b \in \mathbb{T}$. By Lemma 2.1, we have $\operatorname{CPrk}(B) \neq 1$ and so it follows that $\operatorname{CPrk}(G) = \operatorname{cc}(G) = 2$. EXAMPLE 4.4. Let $E_5 = 5K_1$ be an empty graph with 5 vertices and let



be a star graph with six vertices. By Lemma 3.2 and Example 3.7, it follows that

$$\operatorname{CPrk}(S_6) = \operatorname{CPrk}(E_5) = 6 > 5 = \operatorname{cc}(S_6).$$

EXAMPLE 4.5. Let



I L
AS

David Dolžan and Polona Oblak

and assume that $\operatorname{CPrk}(H) = \operatorname{cc}(H) = 2$. Let

$$D = \begin{bmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 2 \\ 0 & 0 & 0 & 2 & 2 \\ 0 & 1 & 2 & 0 & 0 \\ 0 & 2 & 2 & 0 & 0 \end{bmatrix} \in \operatorname{CP}_{5}(\mathbb{T}),$$

and observe that G(D) = H. By Lemma 3.9 (a),

$$D = \begin{bmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 2 \\ 0 & 0 & 0 & 2 & 2 \\ 0 & 1 & 2 & 0 & 0 \\ 0 & 2 & 2 & 0 & 0 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 0 \\ x \\ y \end{bmatrix} \odot \begin{bmatrix} 0 & 0 & 0 & x & y \end{bmatrix} \oplus \begin{bmatrix} 0 \\ w \\ t \\ 0 \\ 0 \end{bmatrix} \odot \begin{bmatrix} 0 & w & t & 0 & 0 \end{bmatrix}.$$

Since $1 = D_{2,4} = \min\{x, w\}$, it follows that x = 1 or w = 1. If x = 1, then $2 = D_{3,4} = \min\{1, t\} \le 1$ and if w = 1, then $2 = D_{2,5} = \min\{y, 1\} \le 1$, both contradictions. Hence, $\operatorname{CPrk}(H) > 2 = \operatorname{cc}(H)$.

Acknowledgment. The authors are grateful to the referees for their helpful remarks and suggestions that improved the presentation of this paper.

REFERENCES

- [1] L.B. Beasley and A.E. Guterman. Rank inequalities over semirings. J. Korean Math. Soc., 42(2):223–241, 2005.
- [2] A. Berman, M. Dür, and N. Shaked-Monderer. Open problems in the theory of completely positive and copositive matrices. Electron. J. Linear Algebra, 29:46–58, 2015.
- [3] A. Berman and N. Shaked-Monderer. Completely Positive Matrices. World Scientific Publishing Co., Inc., River Edge, NJ, 2003.
- [4] I.M. Bomze, W. Schachinger, and R. Ullrich. From seven to eleven: Completely positive matrices with high cp-rank. *Linear Algebra Appl.*, 459:208–221, 2014.
- [5] I.M. Bomze, W. Schachinger, and R. Ullrich. New lower bounds and asymptotics for the cp-rank. SIAM J. Matrix Anal. Appl., 36(1):20–37, 2015.
- [6] D. Cartwright and M. Chan. Three notions of tropical rank for symmetric matrices. Combinatorica, 32(1):55-84, 2012.
- [7] J.H. Drew, C.R. Johnson, and R. Loewy. Completely positive matrices associated with M-matrices. Linear Multilinear Algebra, 37(4):303–310, 1994.
- [8] M. Gondran and M. Minoux. Graphs, Dioids and Semirings: New Models and Algorithms. Operations Research/Computer Science Interfaces Series, Vol. 41, Springer, New York, 2008.
- [9] P. Guillon, Z. Izhakian, J. Mairesse, and G. Merlet. The ultimate rank of tropical matrices. J. Algebra, 437:222-248, 2015.
- [10] L.T. Kou, L.J. Stockmeyer, and C.K. Wong. Covering edges by cliques with regard to keyword conflicts and intersection graphs. Comm. ACM, 21(2):135–139, 1978.
- [11] P. Mohindru. Completely positive matrices over Boolean algebras and their CP-rank. Spec. Matrices, 3:69-81, 2015.
- [12] J. Orlin. Contentment in graph theory: Covering graphs with cliques. Nederl. Akad. Wetensch. Proc. Ser. A 80=Indag. Math., 39(5):406-424, 1977.
- [13] N. Shaked-Monderer. Bounding the CP-rank by graph parameters. Electron. J. Linear Algebra, 28:99–116, 2015.
- [14] Y. Shitov. On the max-min and tropical CP-rank conjectures. Linear Multilinear Algebra, 64(2):219–220, 2016.

